



ERICSSON

THE EVOLUTION OF PROGRAMMABLE NETWORKS: FROM ACTIVE NETWORKS TO SOFTWARE DEFINED NETWORKS (SDN)

Ahmad Rostami, PhD
Ericsson Research, Stockholm
ahmad.rostami@ericsson.com

OUTLINE OF TUTORIAL – I



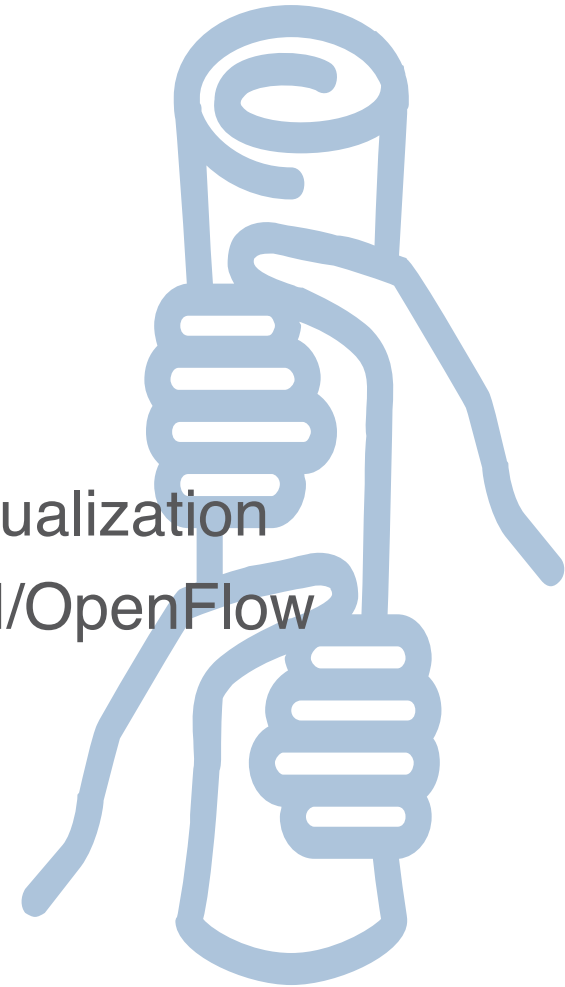
- › Programmability in Telephony Networks
- › Introduction to Programmable Networks
 - › Definition (**What**) & **Why** we need that
 - › How can we evaluate different approaches
 - › Approaches to PN (**How**): A Quick Retrospective Review
 - › Active Networks
 - › (G)MPLS
 - › ForCES



OUTLINE OF TUTORIAL – II



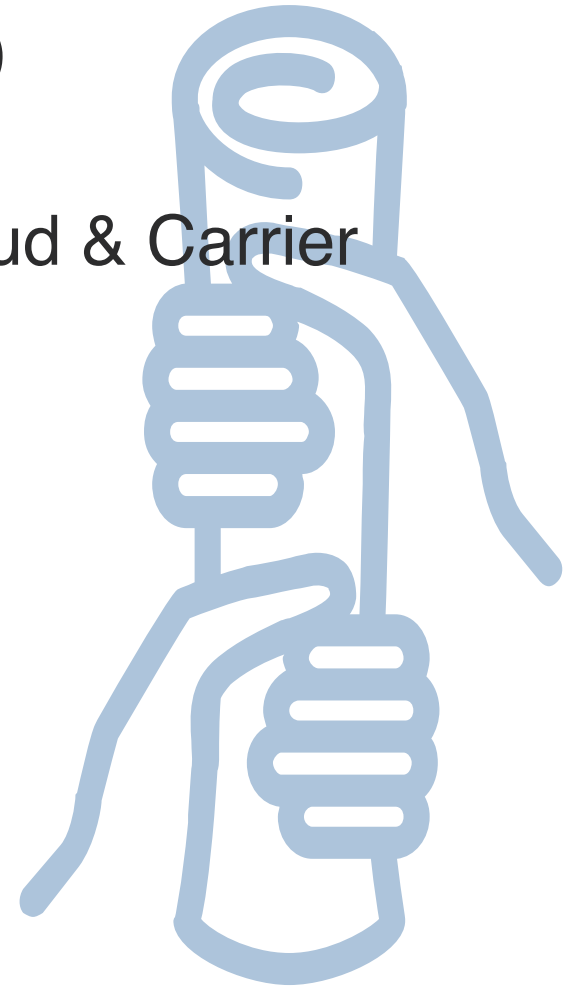
- › Software Defined Networking (SDN)
 - › Introduction & Basic Concept
 - › SDN Realization Example: OpenFlow
 - › SDN/OpenFlow Benefits and Issues
 - › SDN Controllers
 - › SDN Use-case Example: Network Virtualization
 - › Programmability Beyond Current SDN/OpenFlow



OUTLINE OF TUTORIAL – III



- › Network Function Virtualization (NFV)
 - › Relation to SDN
- › SDN and NFV in Action: Unifying Cloud & Carrier Networks
- › Pointers to Relevant Bodies

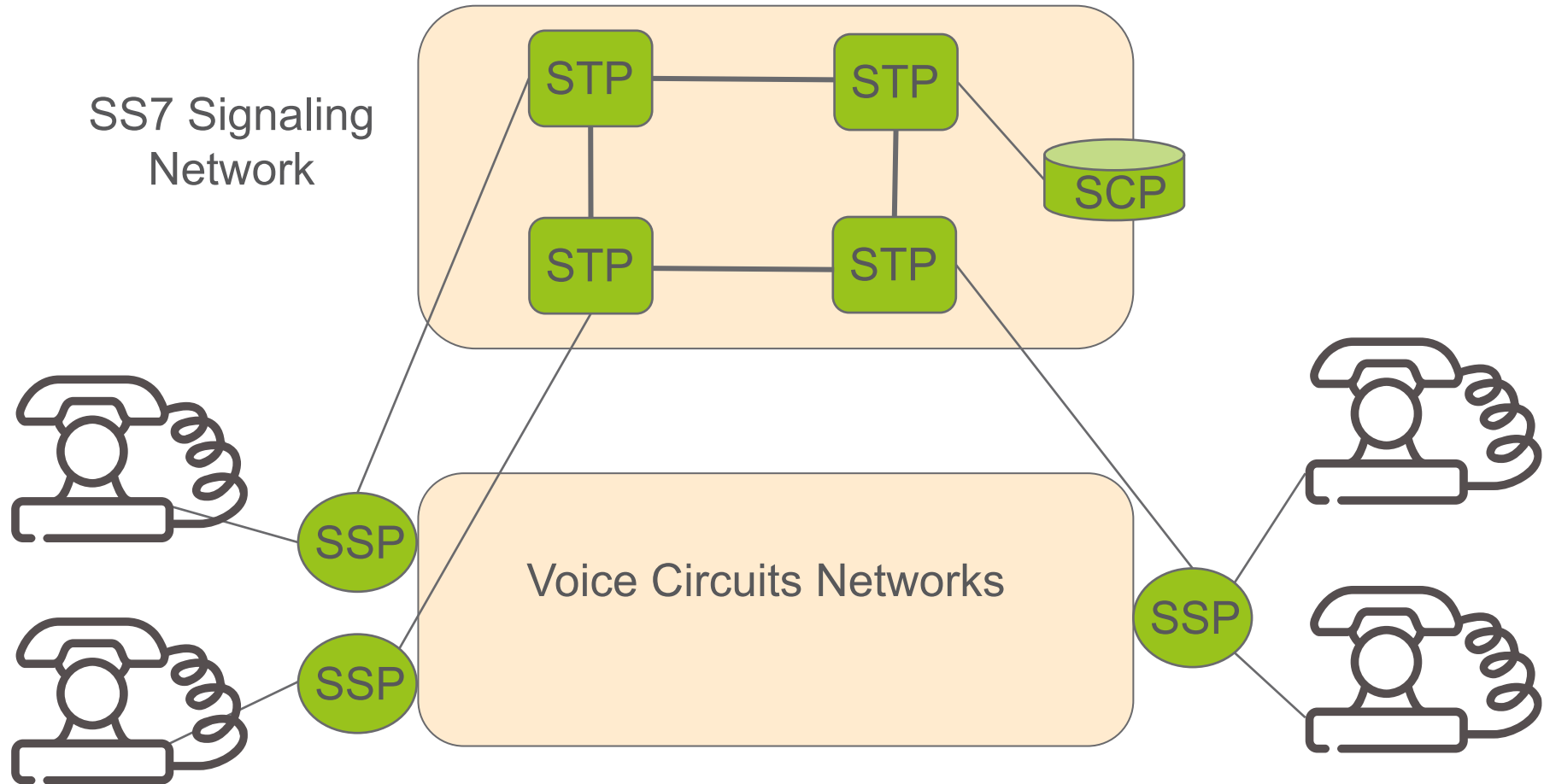


A QUICK LOOK AT THE CONTROL
IN TELEPHONY NETWORKS

—OR—

HOW WAS PROGRAMMABILITY
INTRODUCED INTO TELEPHONY
NETWORKS MORE THAN 20
YEARS AGO?

SIGNALING SYSTEM NO. 7 (SS7)



SSP: Service Switching Point
STP: Signal Transfer Point
SCP: Service Control Point

INTELLIGENT NETWORKS (IN) – I



Value-Added Telephony Services

Introduction of digital switches in 80's opened opportunities for offering new value-added services: call forwarding, call waiting, ...

Approach & Shortcomings

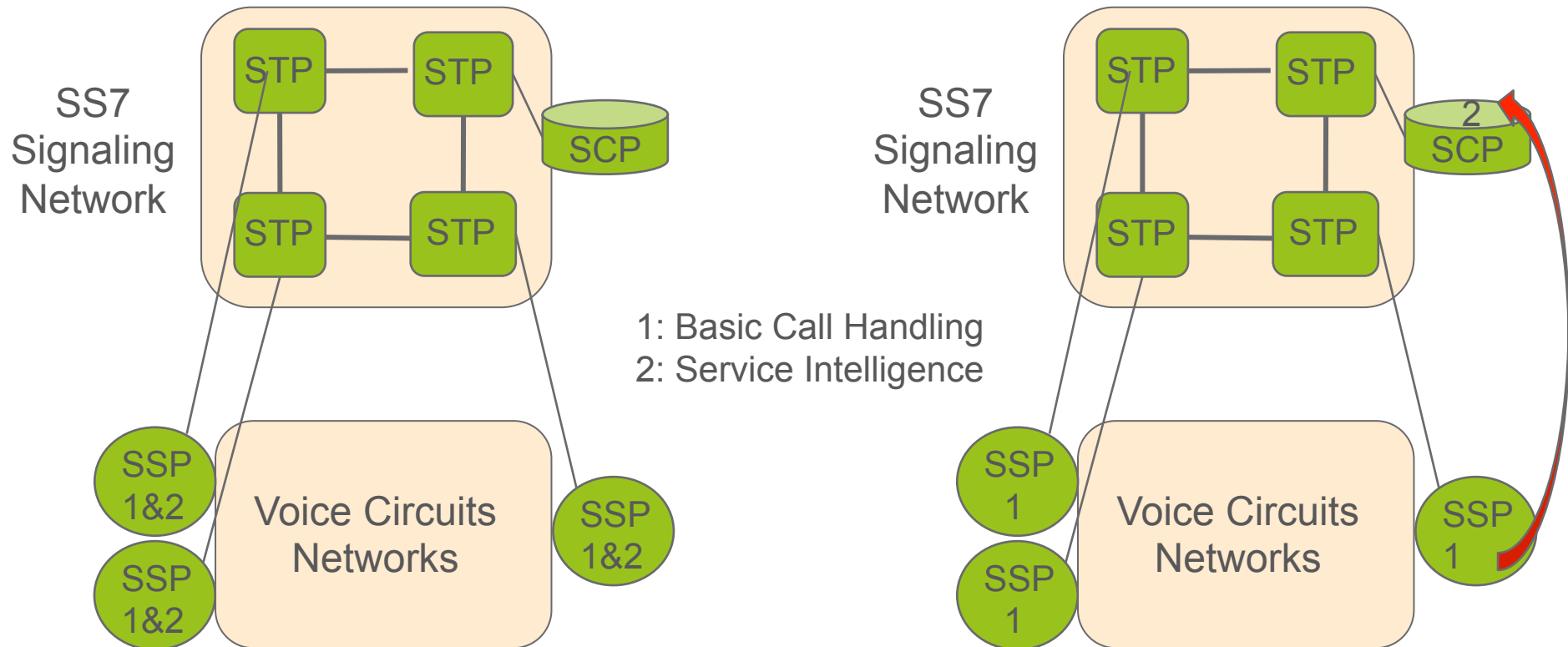
The main approach to new services was based on the pre-programmed switches (SSPs). Lack of Programmability of Switches led to lengthy process of creating new Services.

Introduction of IN

Decouple the service intelligence from the basic call control in a switch.

Push the service intelligence to (a few) centralized location inside the network.

INTELLIGENT NETWORKS (IN) – II



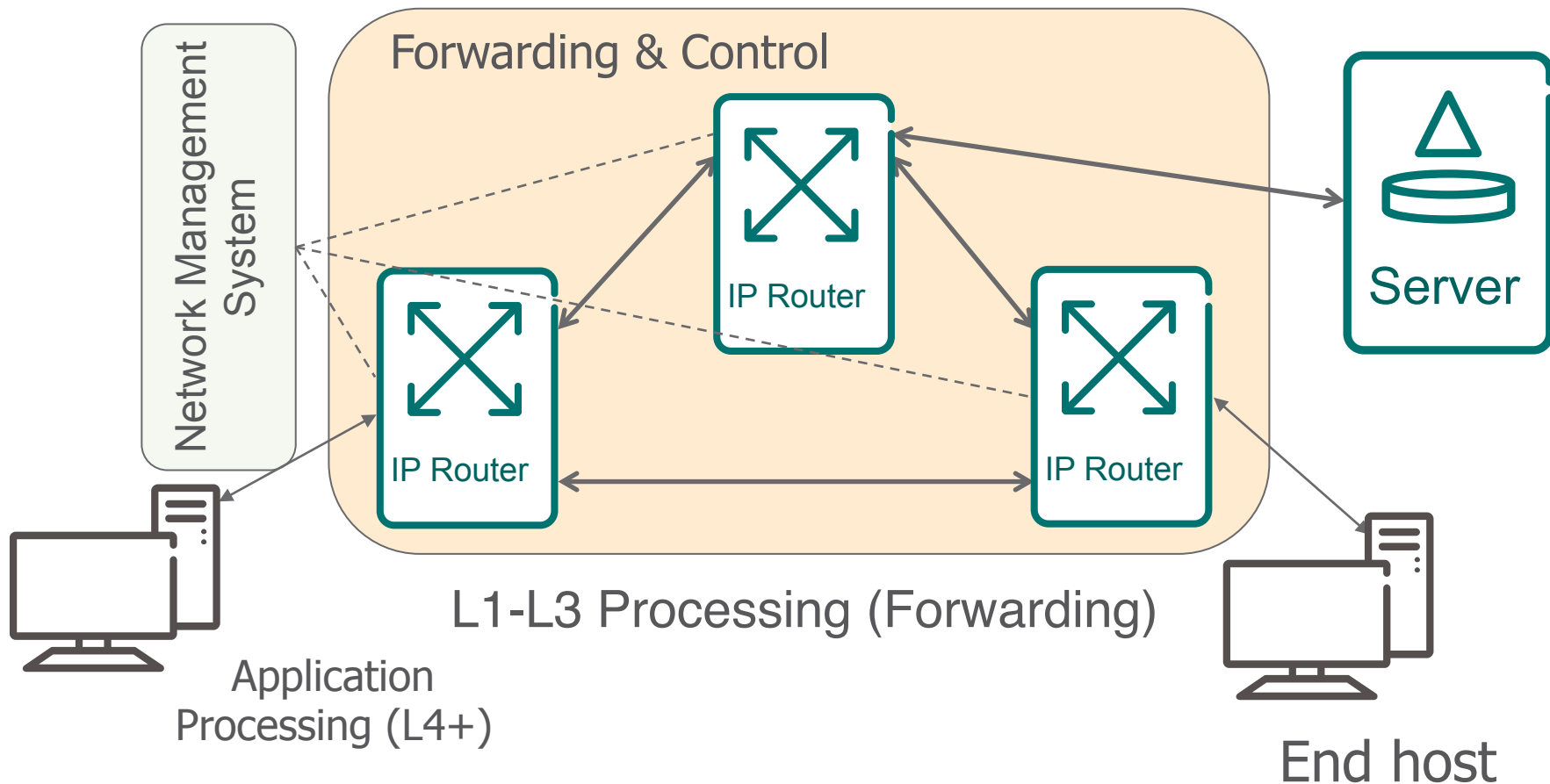
Making the call control programmable enabled (facilitated) many interesting value-added services, e.g., 0800-xxx numbers.

INTRODUCTION TO NETWORK PROGRAMMABILITY

CLASSICAL INTERNET ARCHITECTURE



Internet: an interconnection of comm. nodes each *monolithically* realizing *forwarding & control* based on a *layered architecture*



INTERNET ARCHITECTURE SIMPLICITY



- › Simplicity as a key driving force of the Internet success
 - › Network provides communication between nodes as the best effort service
 - › Keep the network as simple as possible, push the (service) intelligence to the edge

- › BUT, advanced services need more support from the network → Several patches and fixes (look at IETF)

- › Is the Internet architecture still simple?

ISSUES WITH THE INTERNET ARCHITECTURE — I



- › *Complexity of Network Control and Management*
 - › *Layered architecture: blessing and cursing*
 - › *Redundant & contradicting functions at different layers*

- › *Complexity of Changing Network Control*
 - › *Introducing new protocols/protocol stacks is hard*
 - › *Consider how an operator can apply a customized routing mechanism*

ISSUES WITH THE INTERNET ARCHITECTURE — II



› *Complexity of Creating New Services*

- › *Mapping service requirements to vertically integrated devices is difficult*
- › *Applications may need cross-layer interactions*
- › *Applications may need L4+ in-network processing*
- › *Workarounds using inefficient overlays or middle-boxes*

**Need for more flexibility/adaptation in the network
→ Programmable Networks**

WHAT IS NETWORK PROGRAMMABILITY?

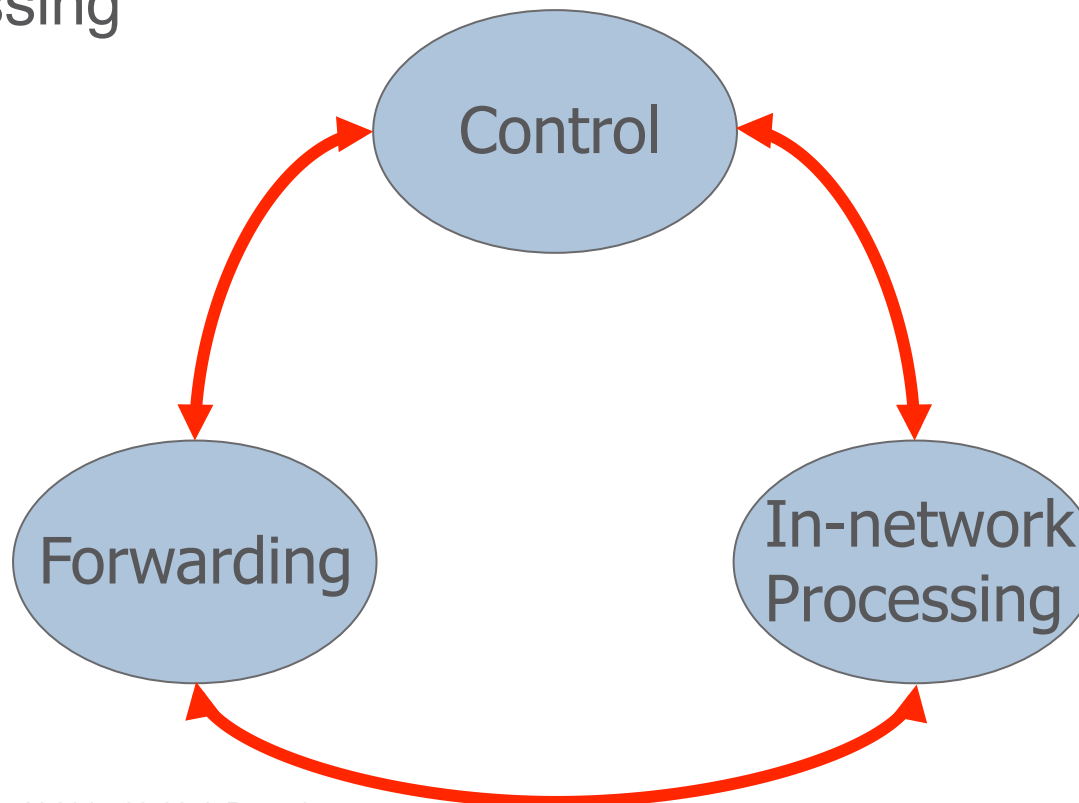


- › The ability to *dynamically* introduce a new architecture, protocols and services on top of a networking infrastructure (without requiring manual configurations)
- › Examples of programmability:
 - › Customize treatments for a group of packets
 - › Assign new semantics to pieces of packets' headers
- › Network Programmability has been a hot topic for research for the last 20+ years, influenced a lot by computing technology

NETWORK PROGRAMMABILITY: HOW? (I)



- › Restructure the main blocks/functions
 - › Where to place the control intelligence?
 - › Balance between basic forwarding and In-network processing



NETWORK PROGRAMMABILITY: HOW? (II)



- › Design (standard) interfaces between the architectural components
- › Select the components that should be open to programming
 - › Programming Interfaces
 - › Programming levels
 - › Who should take control?
- › Important factors: Flexibility, Scalability, Security, Openness, Cost

APPROACHES TO NETWORK PROGRAMMABILITY



Active Networks

- Initiated by academia in mid 90s
- Not deployed in production networks

(G)MPLS*

- Initiated mainly By vendors in late 90s
- Standardization going on in IETF since early 2000s
- Deployed in production networks

ForCES

- Initiated mainly by vendors in early 2000s
- Standardization going on in IETF since mid 2000s
- Not deployed in production networks

SDN

- Initiated mainly by academia in late 2000s
- Standardization going on in ONF since early 2010s
- Partially Deployed in production networks

* (G)MPLS is not designed explicitly as an approach to programmable networks. But, it applies principles which are fundamental to network programmability.

EVALUATION FRAMEWORK



- › What part is programmed?
 - › Forwarding plane, Control plane, ...
- › What is the level of programmability?
 - › Packet, Flow, Class
- › Who takes the control over programmability?
 - › Network Operator, Service Provider, End User
- › Methodology: How is the network programmed?
 - › Separation of control and forwarding plane
 - › Logical vs. physical separation & centralized vs. distributed control
- › Availability and extent of In-network Processing

ACTIVE NETWORKS

ACTIVE NETWORKS – BACKGROUND

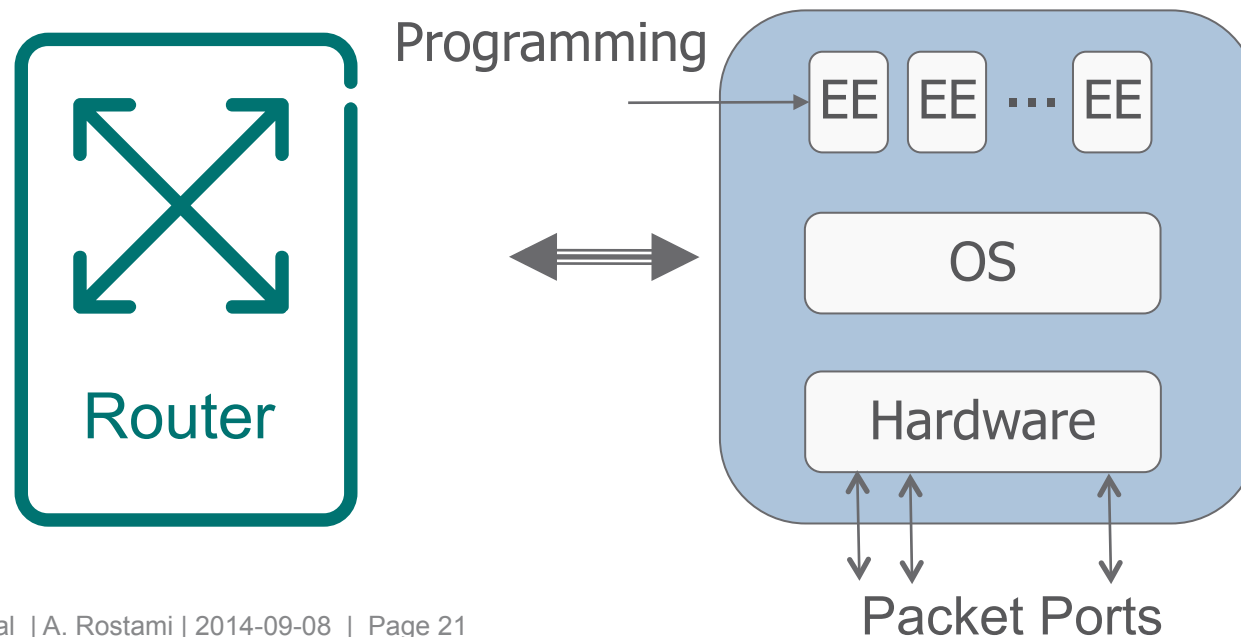


- › Research on Active Networks began in early 90's and supported by several projects
- › Research on Active Networks was motivated by the widespread use of the Internet (applications) & complexity in testing new ideas in real networks
- › To a high extent influenced by the approach in the computing (PC) industry
- › Active Networks requires drastic changes to the existing networking nodes & architecture

ACTIVE NETWORKS PRINCIPLES



- › Replace ad-hoc intelligence in network devices with *generic execution environments (EE)*
- › Expose an interface (e.g. instruction set) to the EE: used to realize a variety of service-specific packet processing



PROGRAMMING THE PROCESSING



- › Service providers &/Or end users (applications) directly *program* network nodes via access to EE to realize desired functionality
- › Customized processing on packets: traffic *actively* select how it is treated/processed

› Programming Approaches:

- › ***Integrated:*** codes are carried together with data packets



- › ***Discrete:*** Code is first loaded into nodes through out-of-band (management) interface and then data packets are injected.

ACTIVE NETWORKS – SUMMARY I



- › Aggressive approach with the highest level of flexibility
- › Programmability directly at the data plane
- › Challenging resource management
 - › Balance between processing and forwarding in nodes
- › Complexity
- › Network operator concerns (e.g., inconsistent applications/codes, security)
- › Lack of a killer application at the time

ACTIVE NETWORKS – SUMMARY II



- › Who Programs the network?
 - › *End user, service provider*
- › What is programmed?
 - › *Network elements, both control and forwarding*
- › Level of programmability
 - › *Packets*
- › Methodology
 - › *No clear separation of control & forwarding*
 - › *Low-level interface to network elements*
- › In-network Processing
 - › *In every node*

Active Networks contributed to the evolution of programmable networks, but as an approach it was not successful.

(GENERALIZED)
MULTI-PROTOCOL
LABEL SWITCHING
(GMPLS)

(G)MPLS – BACKGROUND

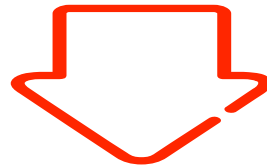


- › Motivation for MPLS introduced in early 2000
 - › Concerns about scalability of IP packet processing in the backbone
 - › Need to attach some notion of connection to traffic
 - › Need for more control over the traffic routing & for traffic engineering
- › Started for packet networks and extended to other switching technologies like TDM and optical circuits → GMPLS
- › A different approach than active networks
 - › Limited flexibility in packet processing

(G)MPLS – PRINCIPLES



- › (Logical) separation of control & forwarding planes
- › Simplify the forwarding plane (in comparison to IP)
- › Push the intelligence/processing from core to edge of networks



- › Partition the entire set of possible packets into a set of Forwarding Equivalence Classes (FECs)
- › Select a path for each FEC → Labeled Switched Path (LSP)
- › At edge of network: process packets and assign them to FECs and label them
- › In core: forward packets only based on their FEC & associated LSPs

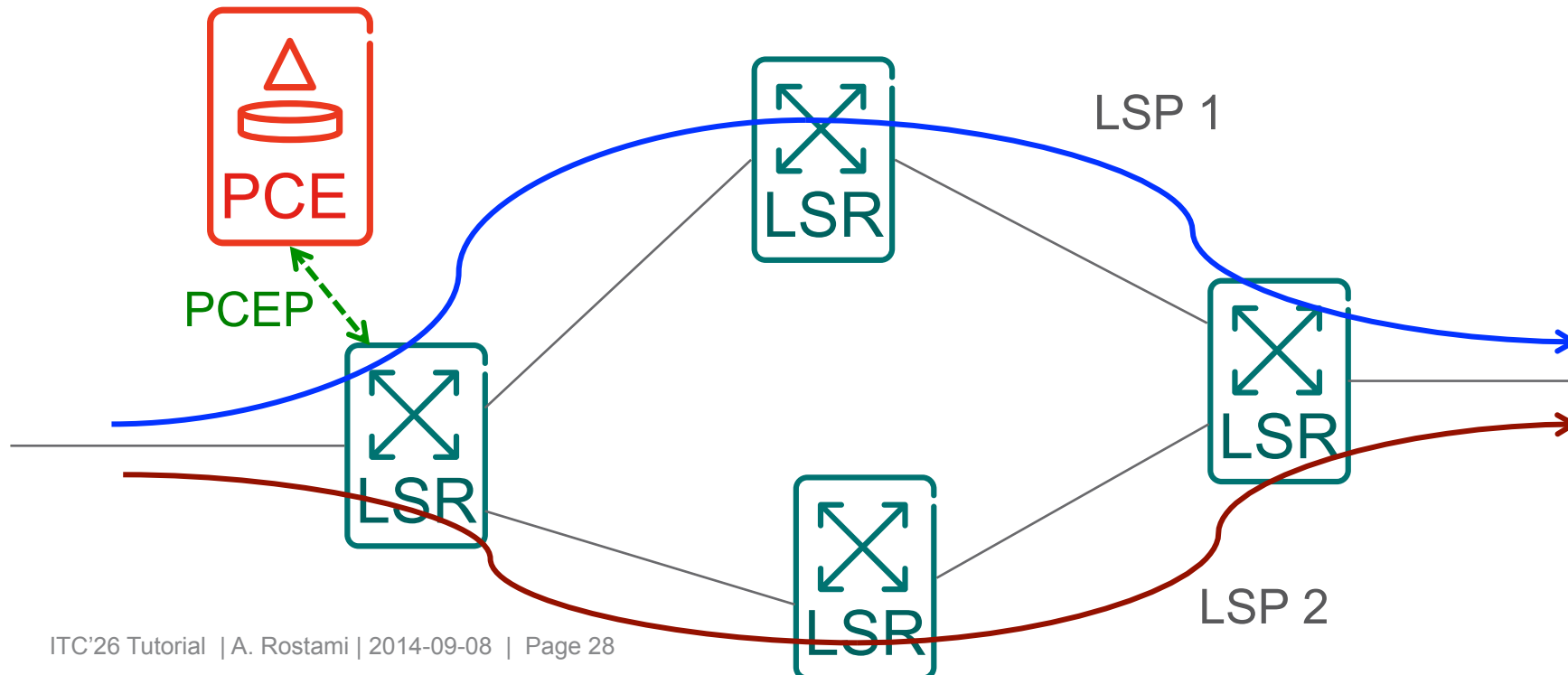
TRAFFIC ENGINEERING IN (G)MPLS NETWORKS



- › Partitioning traffic space into FECs (and LSPs) facilitate more flexible routing and resource allocation in the network

→ constraint-based path selection

- › (Logically) Centralized Path Computation Element (PCE)



(G)MPLS – SUMMARY I



- › Separation of control and forwarding
 - › A step forward towards flexible network control

- › Centralized PCE

- › Programmability limited to constraint-based path computation and
 - › No customized control beyond routing
 - › No in-network processing

(G)MPLS – SUMMARY II



- › Who Programs the Network?
 - › *Network Operator*
- › What is programmed?
 - › *Forwarding treatment, path selection*
- › Level of programmability
 - › *Traffic classes*
- › Methodology
 - › *Logical separation of control & forwarding*
- › In-network Processing
 - › *Not available*

FORWARDING &
CONTROL ELEMENT
SEPARATION
(FORCES)

FORCES – BACKGROUND



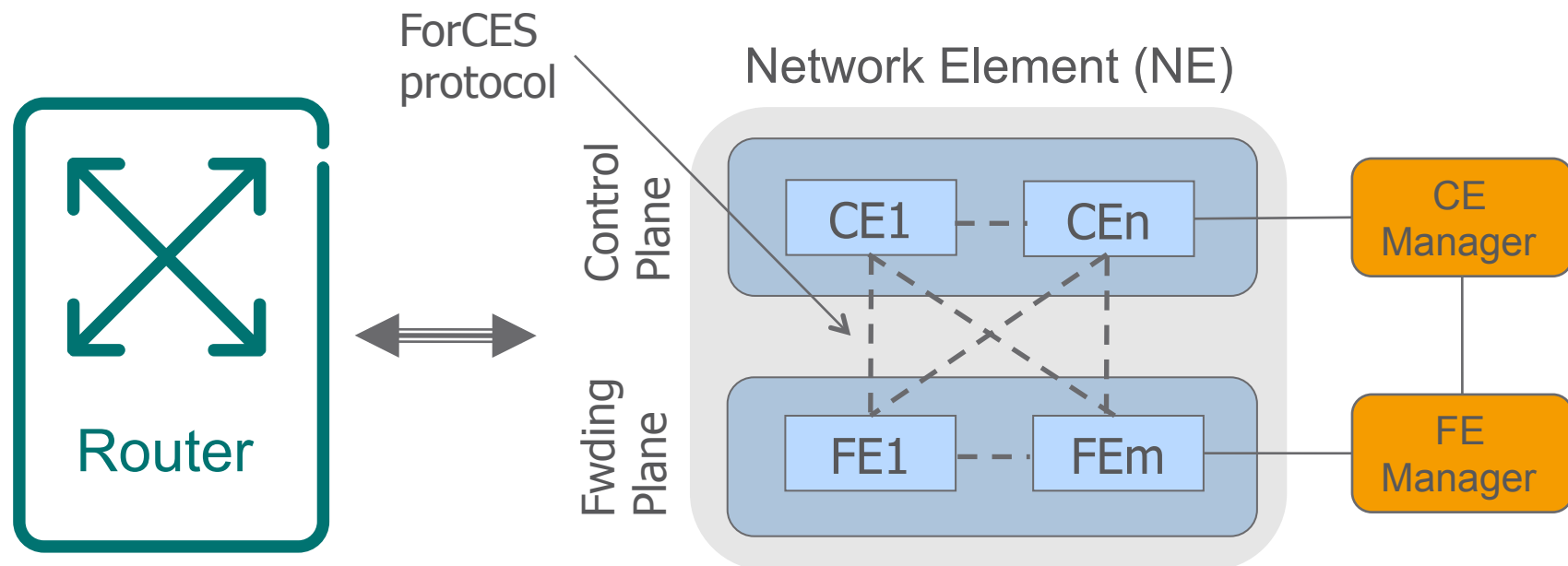
- › Motivations for ForCES introduced in early 2000
 - › Separating Control & Forwarding in network elements
 - › Powerful Network Processing Units (NPUs) allowing realization of customized functions in data plane
 - › Need for a modular functional architecture describing components in data and control planes
 - › Need to specify (& standardize) the interface between components of the functional architecture

- › A Modular architecture supporting network programmability
 - › A great support for in-network processing
 - › Object-oriented programming concept applied to the nodes' components management

FORCES – PRINCIPLES

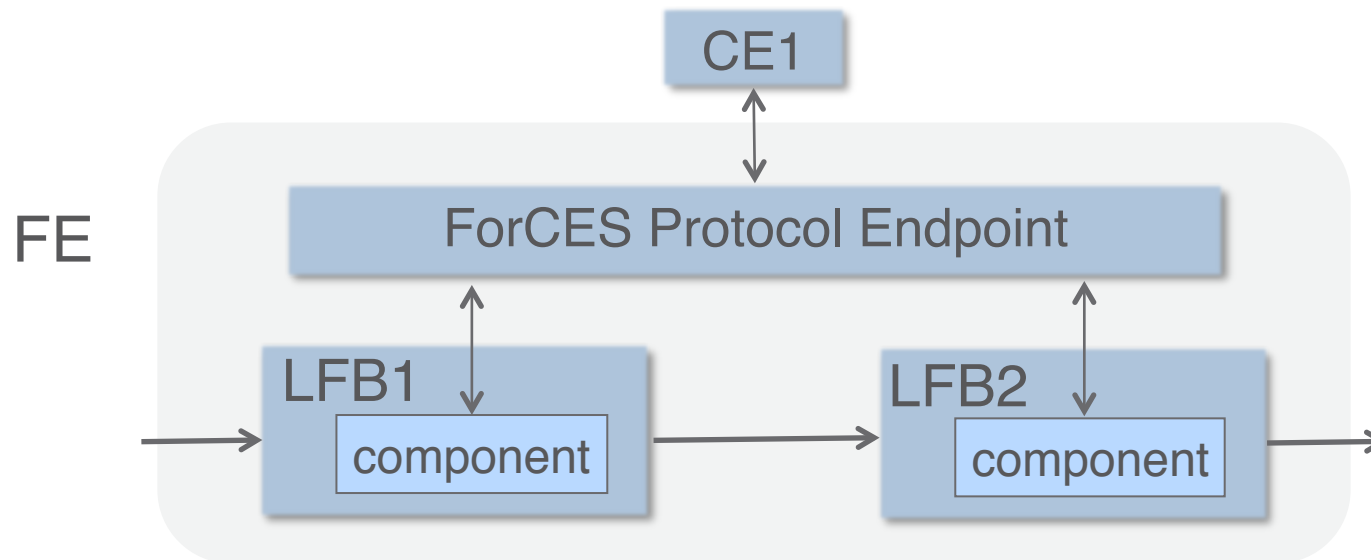


- › (Logical) separation of control & forwarding planes
- › Modular (& dynamic) interconnection of architectural elements



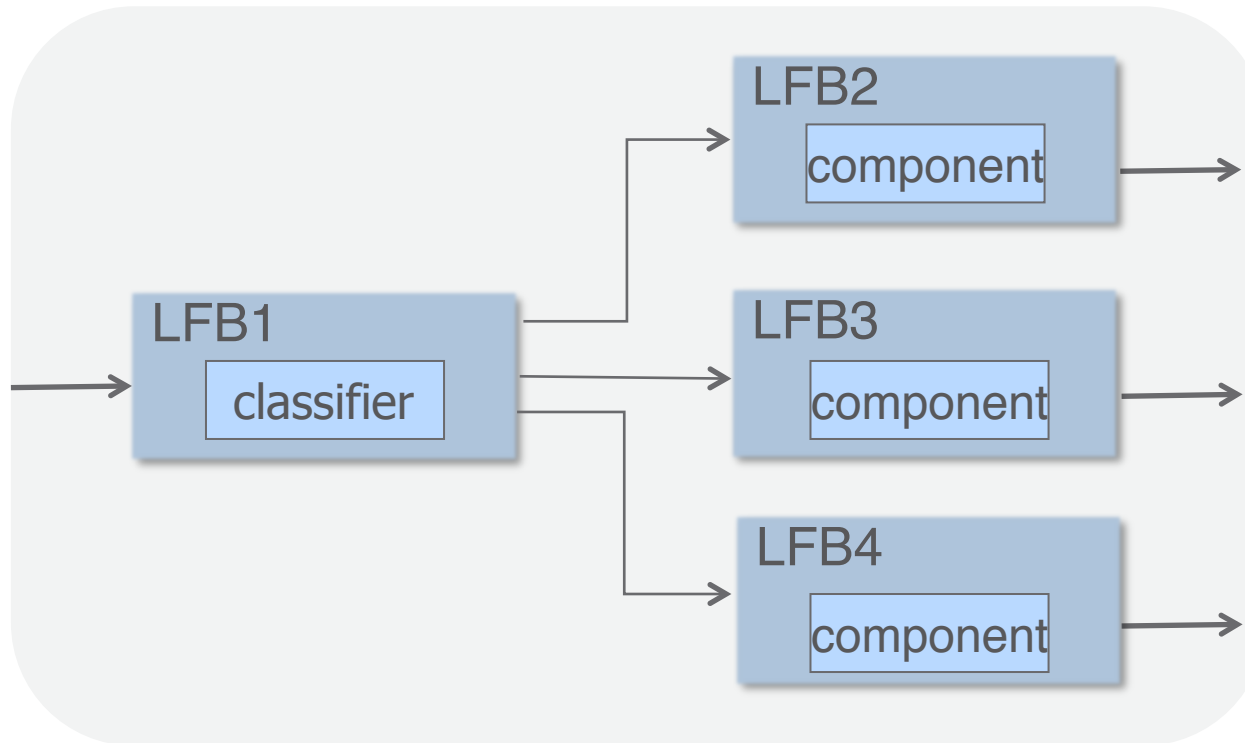
- › Multiple instances of control elements (CE) & forwarding elements (FE) in an NE → interconnected over variety of technologies

PROGRAMMABILITY WITH FORCES



- › Forwarding model is based on an abstraction of logical functional blocks (LFB)
- › Each LFB defines a single action on passing packets
- › FE is composed of multiple LFBs interconnected in a directed graph
- › A CE can *dynamically* create, instantiate, update or delete LFBs within a FE from LFB library

SERVICE COMPOSITION WITH FORCES



- › The outcome of a LFB processing determines how a packet will be processed in downstream LFBs
→ Flow/Class-based Programming

FORCES – SUMMARY I



- › Standard interface between control and forwarding
 - › Independent evolution of the two components
- › More flexibility & generality than MPLS
 - › MPLS functions can be implemented in ForCES
- › Supports a gradual deployment of ForCES nodes
 - › Transparent interconnection to conventional nodes
- › Support for in-network processing
- › No programming interface to control plane

FORCES – SUMMARY II



- › Who Programs the Network?
 - › *Network operator*
- › What is programmed?
 - › ***Forwarding & control***
- › Level of programmability
 - › *Class/Flow*
- › Methodology
 - › *Logical separation of control & forwarding*
 - › *Modular functional architecture*
 - › *Standard interface between modules*
- › In-network processing
 - › *Possible at every node*

SOFTWARE DEFINED NETWORKING (SDN)

SDN – BACKGROUND



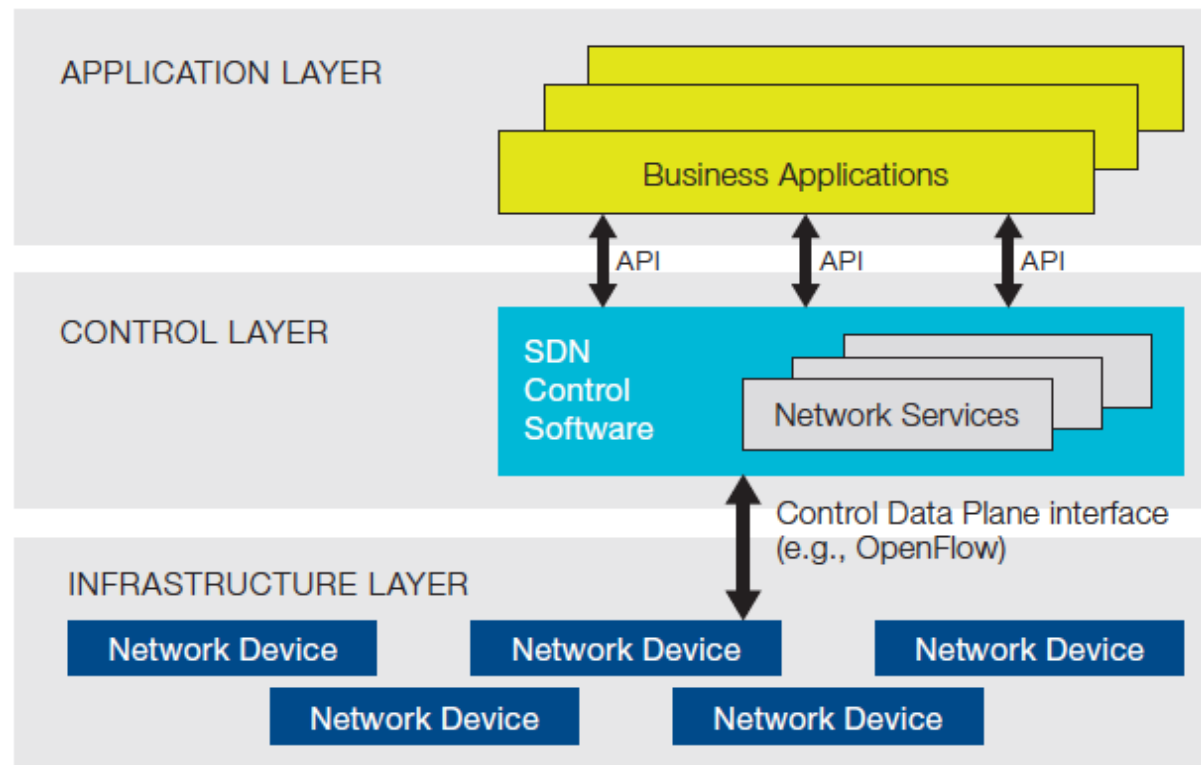
- › Motivations for SDN introduced in late 2000
 - › Separating control and forwarding (!)
 - › No success with previous attempts to programmable networks
 - › The need for (large scale) experimental networking platforms

- › SDN borrows several concepts from previous approaches to programmable networks & computing, e.g.,
 - › Separation of control from forwarding
 - › High level programming interfaces

SDN – PRINCIPLES



- › Physical separation of control & forwarding
 - › Specify an (open standard) interface between control and forwarding
- › Delegate control to a logically centralized controller
- › Abstract the network resources from higher layer applications

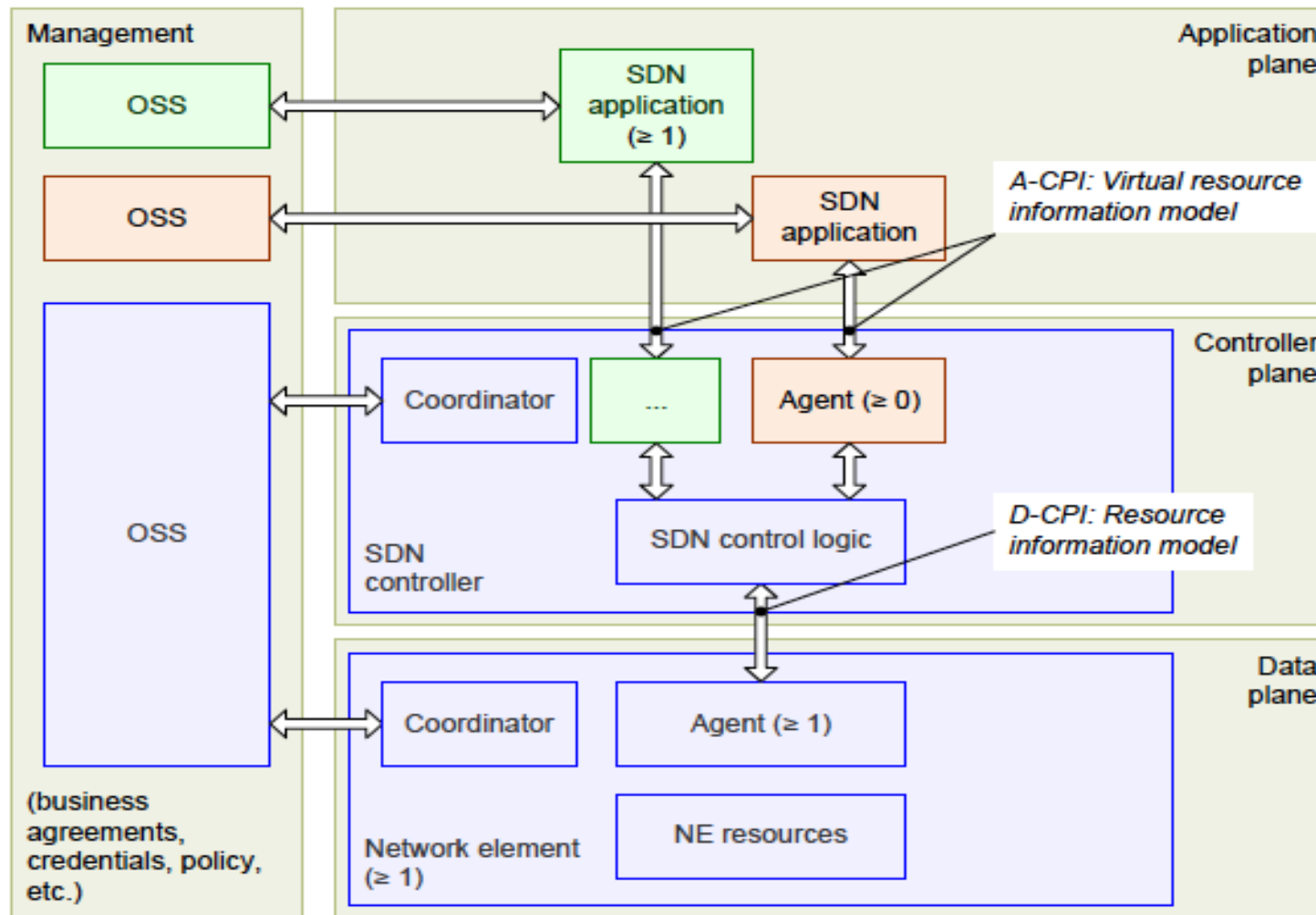


SDN ABSTRACTION LAYERS



- › SDN makes use of abstraction layers to simplify the programmability
- › Forwarding abstraction
 - › Design a Control-Forwarding interface & centralize the control logic
- › Control-plane abstractions
 - › Hide details of network resources & state distribution from control applications
- › There are more than one way to do the abstractions
- › Abstractions can be designed to allow a recursive control architecture

SDN ARCHITECTURE – ONF VIEW



BENEFITS OF SDN – I



- › Control plane can be directly programmed
 - › Customized control logics
 - › Rapid creation of new services
- › Simplified programming of new applications
 - › Programming on an abstract view of network
 - › State distribution carried by the controller (Network OS)
- › Enabling more efficient control applications
 - › Controller provides a global view of resources

BENEFITS OF SDN – II



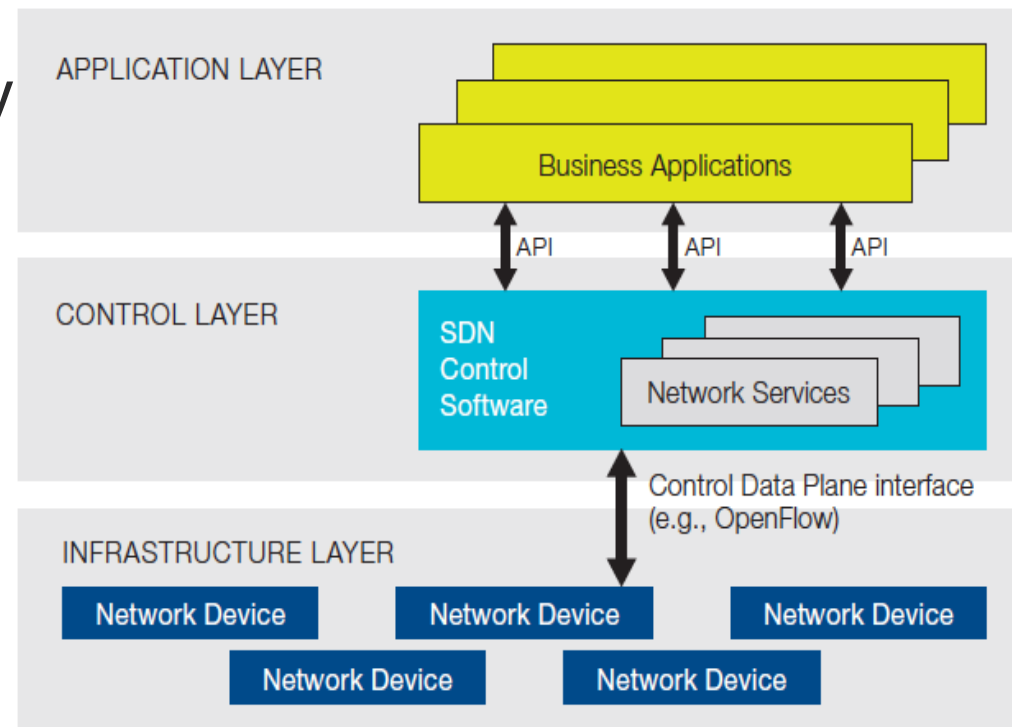
- › Vendor independent network control
 - › Using a standard API between forwarding & control
- › Simplified forwarding elements
 - › Support only a generic API instead of many protocols
- › Data & control planes can evolve independently
 - › Higher rates of innovations
 - › More competitive market

SDN REALIZATION
EXAMPLE:
OPENFLOW

SDN REALIZATION



- › SDN is an architectural framework for programmability
- › A realization of SDN requires concrete:
 - › Data-plane abstraction & API
 - › SDN Controller (Network OS)
 - › Network abstraction and Control API



OPENFLOW — A CONCRETE REALIZATION OF SDN

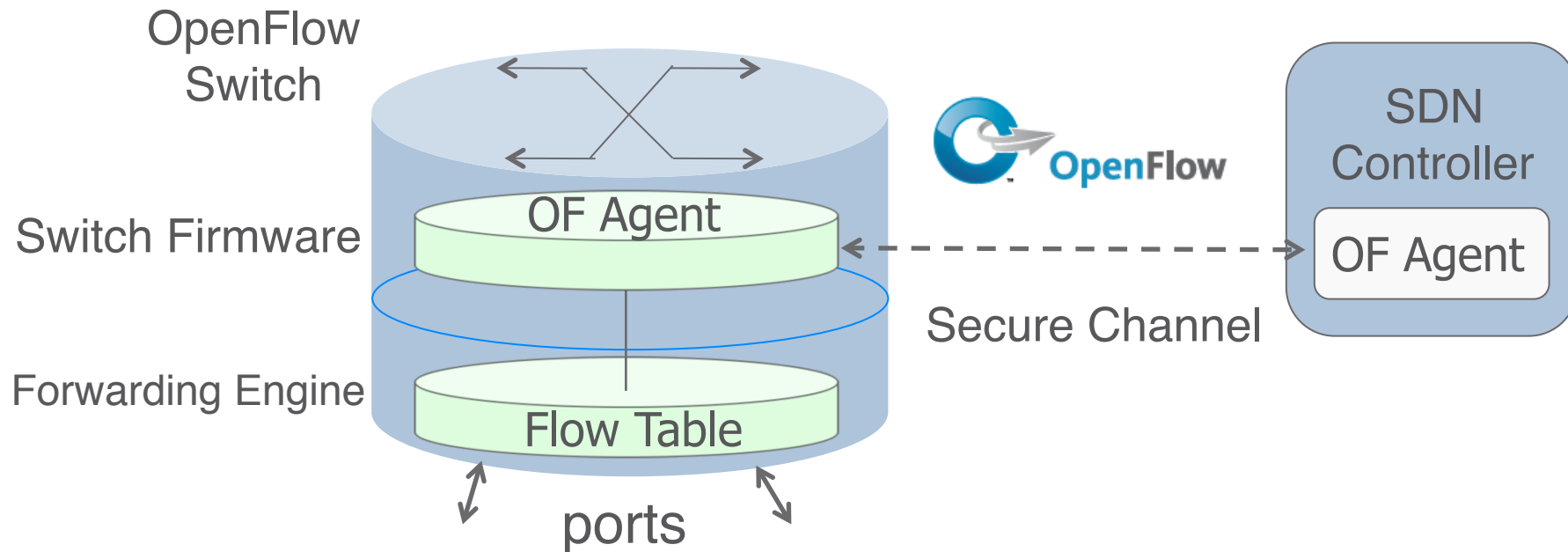


- › OpenFlow is one of the initial works leading to SDN
 - › Started in academia in late 2000
 - › Initially focused on utilizing Ethernet-based campus networks as an experimental platform
- › OpenFlow specifies a forwarding plane abstraction together with an API to forwarding devices
- › OpenFlow does not specify the network controller, nor does it specify a controller-application API

FORWARDING ABSTRACTION



- › OpenFlow started with abstracting data plane of the Ethernet switch
- › In OpenFlow a forwarding element is abstracted as a flow table applying the **Match+Action** principle



- › OpenFlow specifies an instruction set to manage the flow table from a remote controller

FLOW TABLE



Match	Priority	Counters	Instructions	Timeouts	Cookie
-------	----------	----------	--------------	----------	--------

- › Match: A flow is flexibly defined as a group of packets sharing certain header values. Match + Priority → unique flow identifier

Ingress Port	Ether Src	Ether Dst	Ether Type	VLAN	IP Src	IP Dst	IP Proto	L4 Src	L4 Dst
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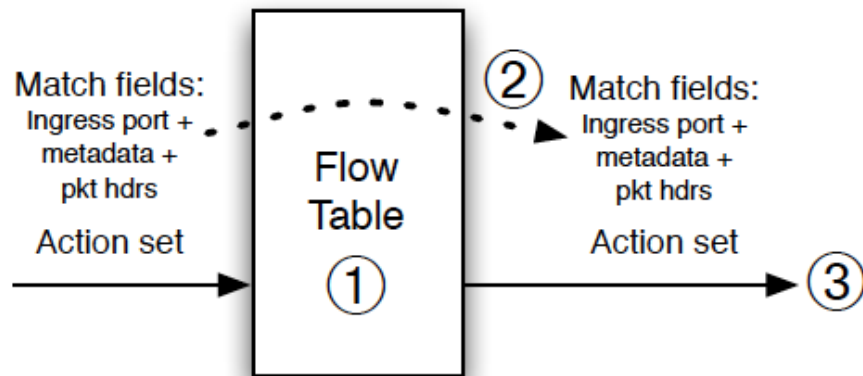
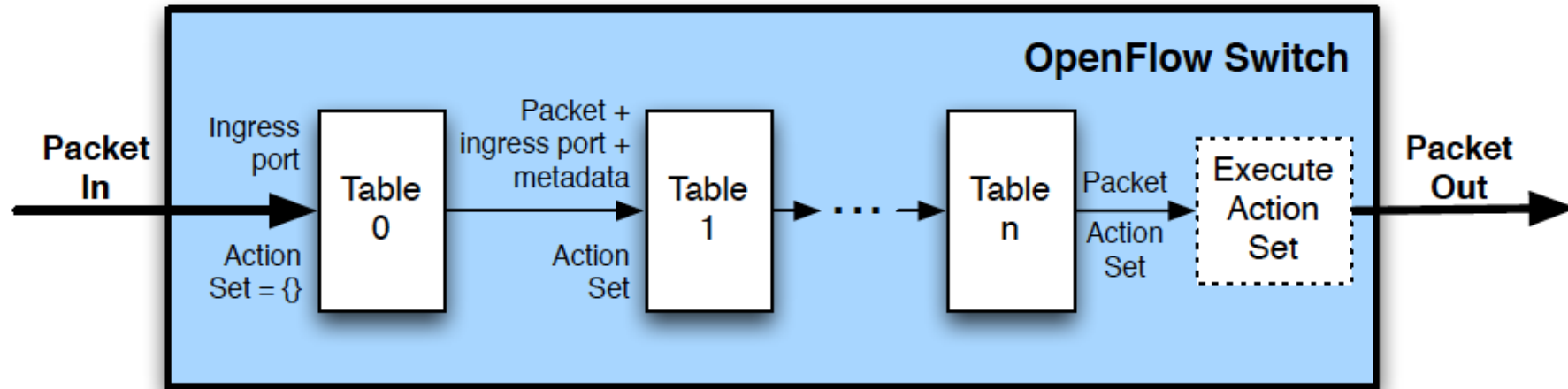
- › Instructions: modify actions set associated with a flow: sent to a port, modify a field, etc.
- › Counters: collect statistics
- › Timeout: specify validity of a flow entry in time
- › Cookie: used by controller for filtering groups of flows

FLOW TABLE: AN EXAMPLE



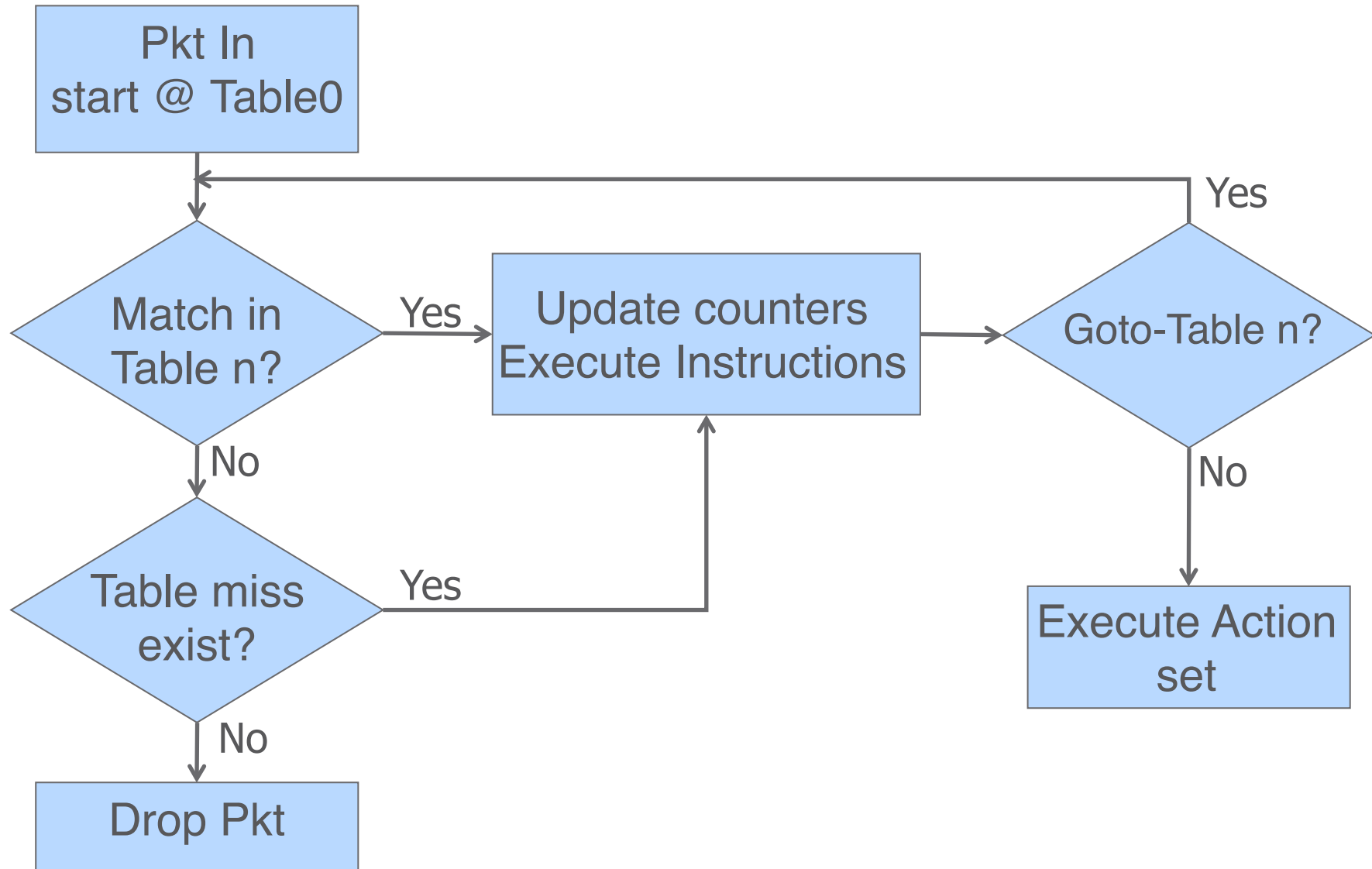
Ingress Port	MAC Src	MAC Dst	Ether Type	VLAN	IP Src	IP Dst	IP Proto	L4 Src	L4 Dst	Actions
*	*	00...	*	*	*	*	*	*	*	Port 2
3	*	*	*	*	*	*	*	*	22	Drop
1	*	*	*	*	*	192. 168. 1.0/2 4	*	*	*	Port 5
*	*	*	*	*	*	*	*	*	*	Send to Controller

DATA-PATH PROCESSING PIPELINE



- ① Find highest-priority matching flow entry
- ② Apply instructions:
 - i. Modify packet & update match fields
 - ii. Update action set
 - iii. Update metadata
- ③ Send meta data and action set to next table

PACKET FLOW IN SWITCH



SDN/OPENFLOW BENEFITS



- › Flow-level programming: OpenFlow enables flexible definition of flows & assignment of individual forwarding treatment to them
- › Flat Control: OpenFlow collapses L1-L4 control into a single layer → multi-layer optimization
- › Traffic Statistics per flow, table, port, ...



- › OpenFlow provides a powerful tool for efficient control of enterprise networks
- › SDN/OpenFlow already applied to WANs & large datacenters

SDN/OPENFLOW ISSUES/OPEN PROBLEMS – I



- › Programmability of data-path packet processing limited by OpenFlow API
 - › No direct support for stateful processing in data-path
- › Protocol dependent nature of OF API
 - › Matching fields extended from 12 in OF 1.0 to 41 in OF 1.4
- › No network abstraction atop OF API
 - › Developing control modules still require dealing with low-level flow API

SDN/OPENFLOW ISSUES/OPEN PROBLEMS – II



› Scalability

- › Controller processing, flow-table size, OF channel, ...

› Interface among Controllers

- › SDN/OpenFlow targets single domain
- › Multi-domain networks, distributed control

› Security concerns

SDN/OPENFLOW SUMMARY



- › Who Programs the Network:
 - › Service provider
- › What is programmed?
 - › (mainly) control plane
- › Level of programmability:
 - › Traffic flow
- › Methodology:
 - › decoupling control intelligence from datapath
 - › Abstract forwarding based on “match+action” flow tables
 - › An open interface between controller & forwarding
- › In-network processing:
 - › Not supported directly

SDN CONTROLLER

SDN CONTROLLER



- › SDN Controller acts as a network operating system
 - › Abstracting the network resources
 - › Providing a programmatic interface for developing network applications in SW
 - › Configuring networking resources as a service to network applications (state distribution)
 - › Ensuring consistent usage of resources among control Apps
 - › Enabling virtualization of networking resources

SDN CONTROLLER PLATFORMS

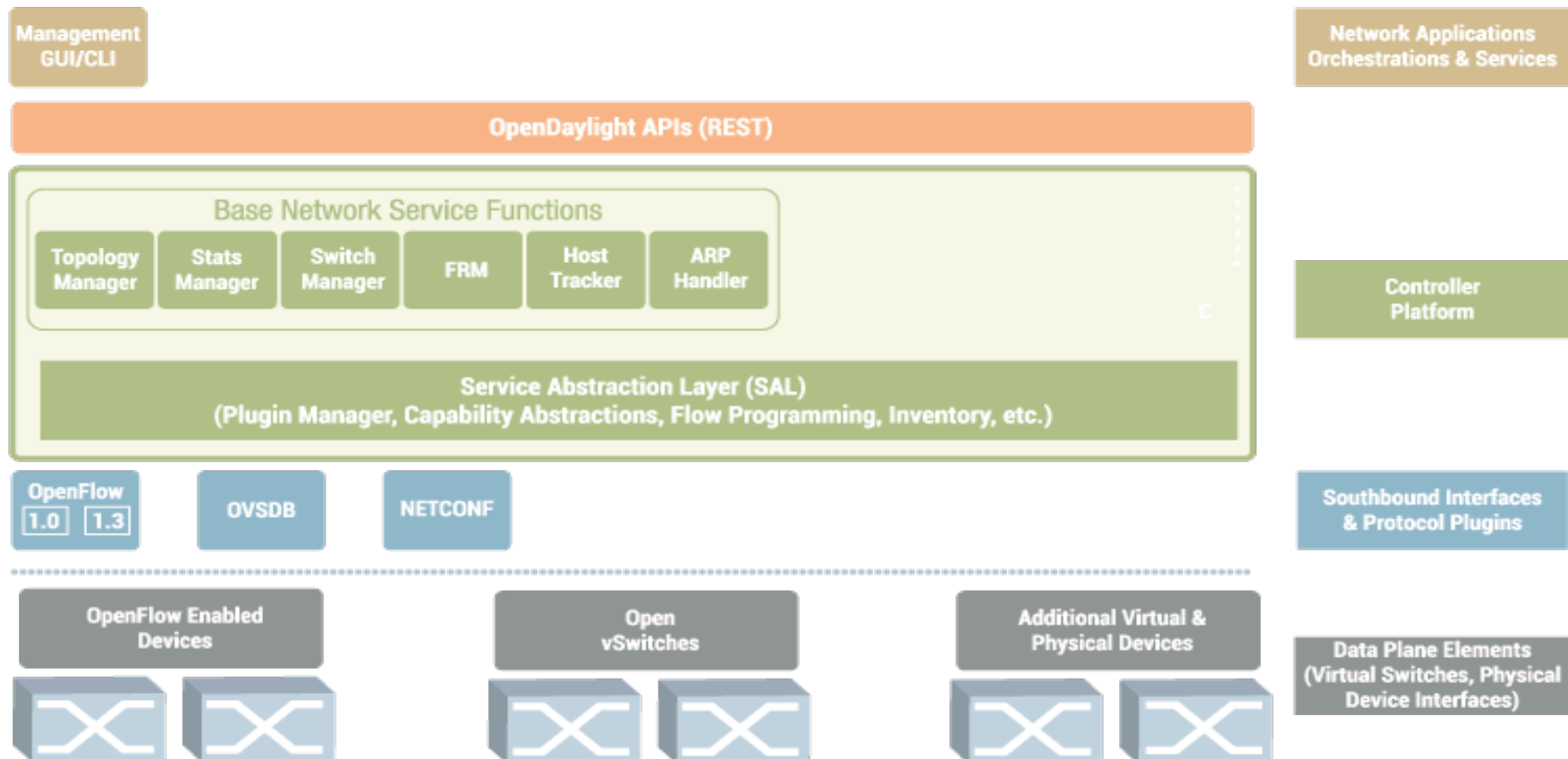


- › Several commercial and educational SDN controllers developed in recent years, primarily focusing on OpenFlow
- › **OpenDayLight (ODL):** An open source collaborative project started in 2013 to develop a common platform for SDN controller
- › ODL supports several control-dataplane interfaces
- › ODL project runs under the umbrella of Linux Foundation, with contribution from many companies:
 - › Ericsson, IBM, Cisco, HP, Microsoft, Juniper, ...
- › First ODL implementation released Feb. 2014 with code name Hydrogen

OPENDAYLIGHT ARCHITECTURE



- VTN:** Virtual Tenant Network
- oDMC:** Open Dove Management Console
- D4A:** Defense4All Protection
- LISP:** Locator/Identifier Separation Protocol
- OVSDB:** Open vSwitch DataBase Protocol
- BGP:** Border Gateway Protocol
- PCEP:** Path Computation Element Communication Protocol
- SNMP:** Simple Network Management Protocol
- FRM:** Forwarding Rules Manager
- ARP:** Address Resolution Protocol



OPENDAYLIGHT SAL



- › SAL is the a major architectural component of ODL, differentiating it from conventional SDN controllers
- › SAL abstracts services and capabilities of various data-plane elements towards higher functional layers of controller

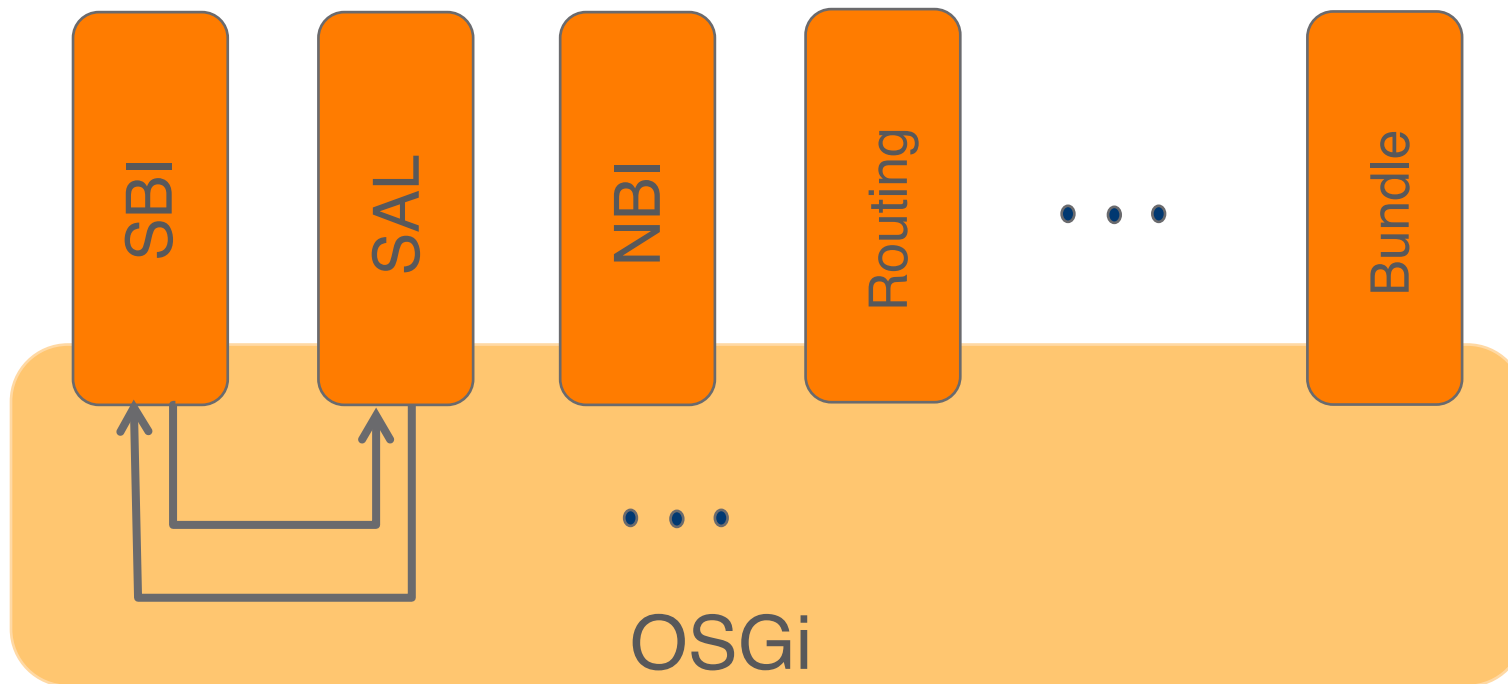
Major SAL Services:

- › **DataPacket:** dispatching data-packets towards higher layers for processing
- › **Flowprogrammer:** allows higher-layer modules to request data-plane programming at the flow level
- › **Topology:** allows to collect topology from different SBI plugins and to stich them in higher functional modules
- › **Inventory:** provides inventory data to applications, etc.

OPENDAYLIGHT SW ARCHITECTURE



- › ODL is mainly programmed in JAVA
- › The Open Service Gateway initiative (OSGi) framework is used to support the modularity



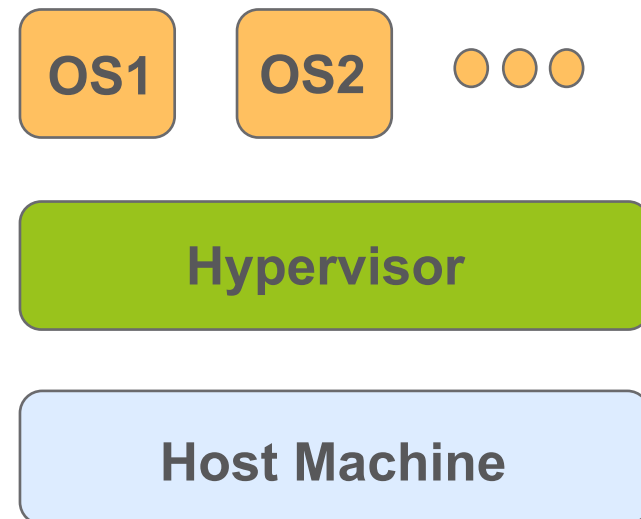
SDN USE-CASE
EXAMPLE:
NETWORK
VIRTUALIZATION

RESOURCE VIRTUALIZATION

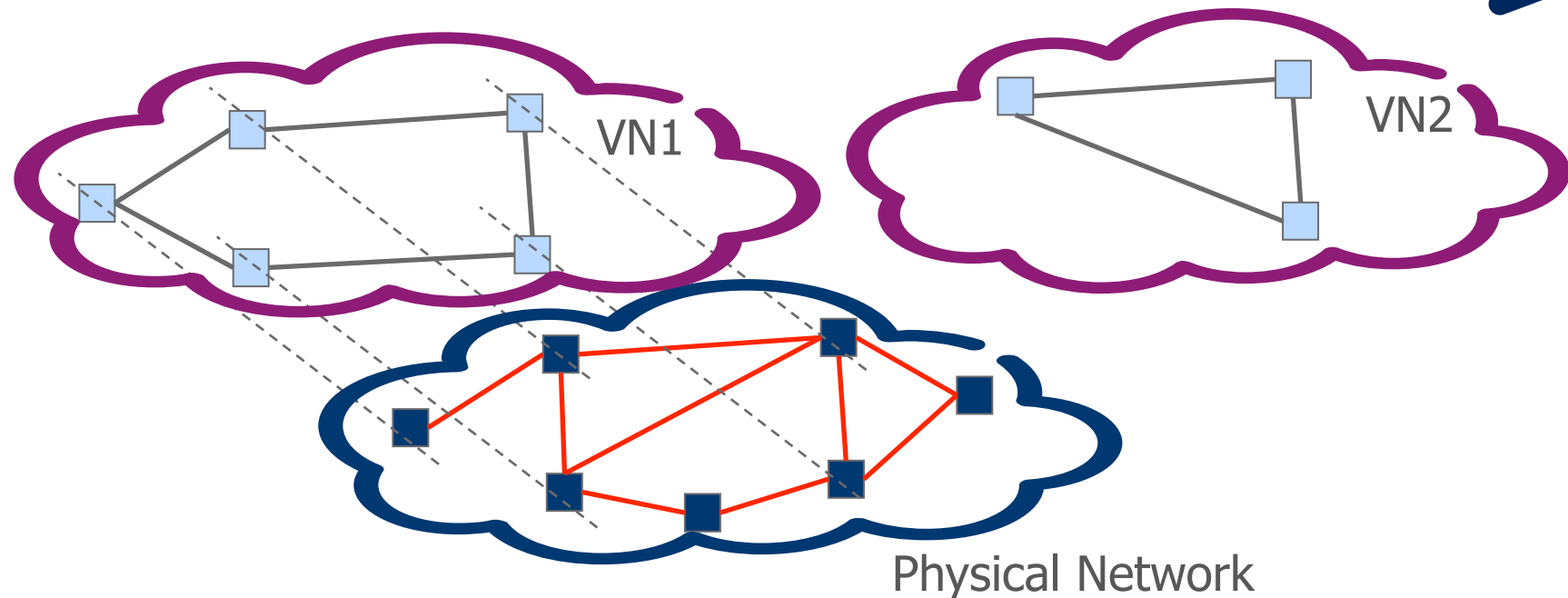


- › Virtualization: creation of (multiple) logical instances of a physical entity (element/resource/functionality)
 - › logical entities share the resources available in physical one
 - › each logical entity can be controlled/managed independently

- › Server Virtualization is a common practice in IT industry:
 - A hypervisor allows several OSes to share a single host machine.



NETWORK VIRTUALIZATION



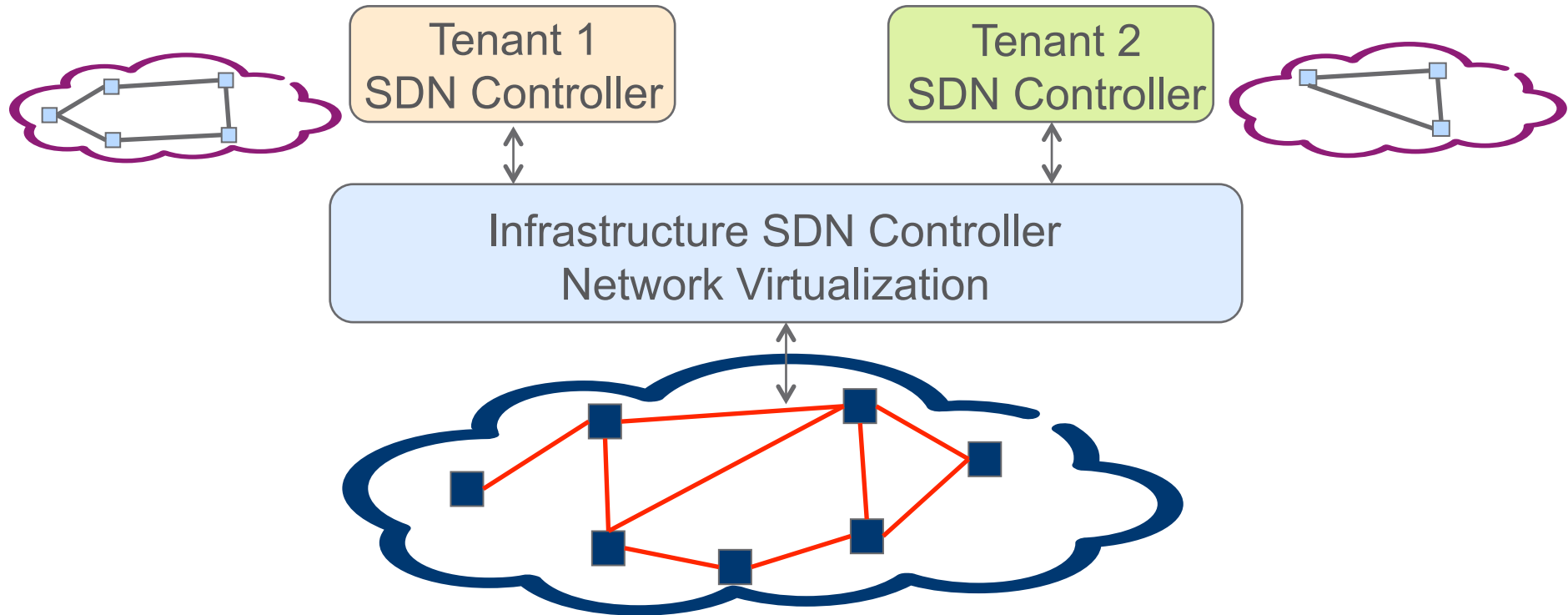
Resource Isolation between virtual networks & Independent Control

› Network Virtualization

- › Link Virtualization: slice/aggregate link capacity
- › Node Virtualization: slice/aggregate packet forwarding and processing resources within switches/routers

› Several approaches to NV: e.g., network overlay, SDN

NETWORK VIRTUALIZATION WITH SDN



- › Infrastructure controller manages the slicing & ensures consistent resource usage between the tenants
- › Each tenant network has full control over its (virtual) network, e.g., for routing
- › Helpful in datacenter environment together with server virtualization

PROGRAMMABILITY
BEYOND CURRENT
SDN/OPENFLOW

PROGRAMMABILITY BEYOND CURRENT SDN/OPENFLOW



- › SDN approaches investigated so far (e.g. OF) merely expose more control knobs for programming the network
- › The knobs' functions are dictated by the fixed functionality of forwarding devices
- › What if we need to (dynamically) reprogram the functionality of forwarding devices?
 - › E.g., to define new protocol stacks and associated packet parsing and processing

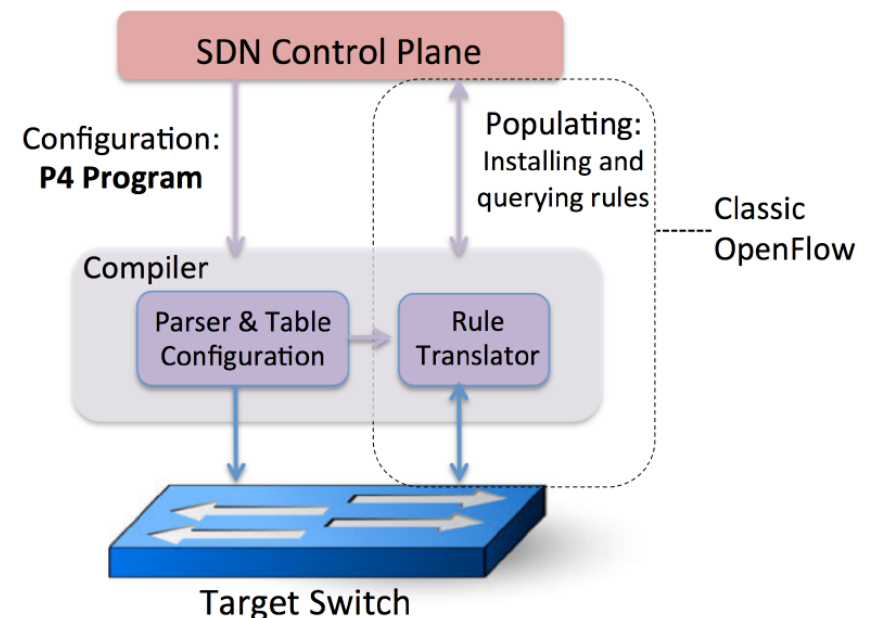
PROTOCOL INDEPENDENT PACKET FORWARDING



- › Protocol Independent Forwarding adds another dimension to the programmable networks.
 - › (Re)configure the forwarding element → specify packet parsing & processing model
 - › Control the forwarding operation (traffic dependent)

› Application to SDN/OpenFlow

- › Configure the parsing & forwarding models in “match +action” tables
- › Populate the tables and control the traffic forwarding



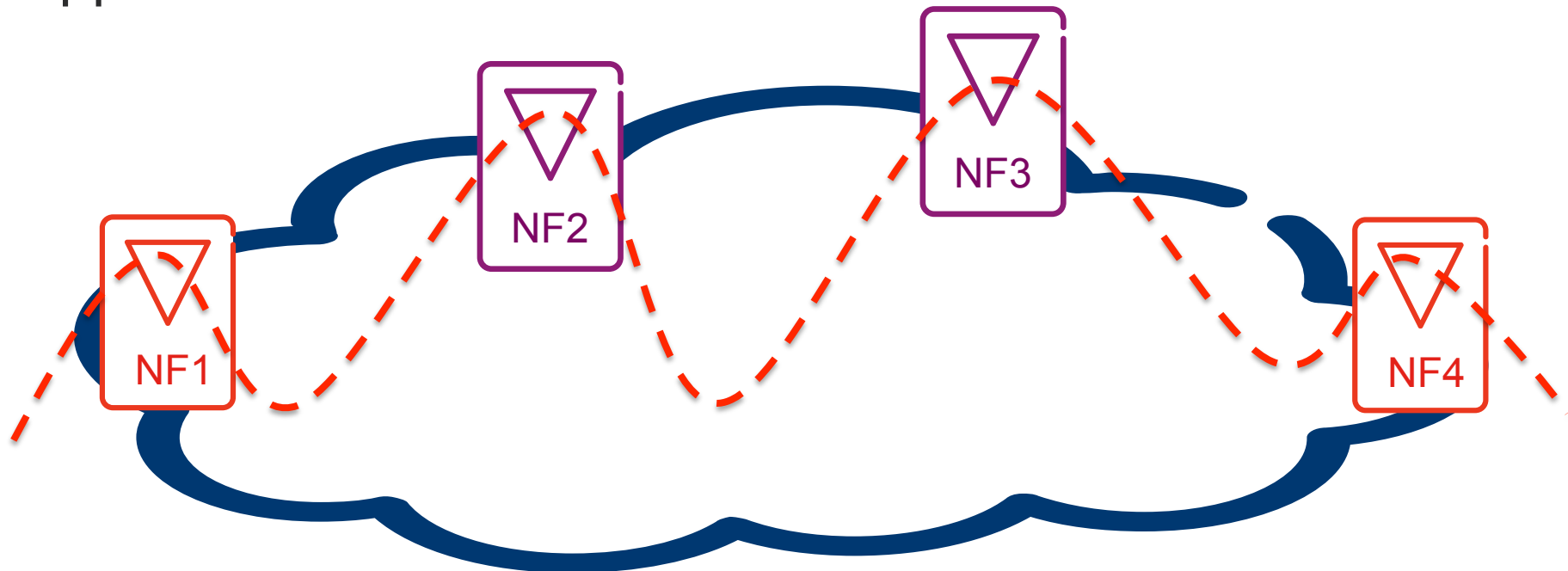
Source of Fig.: P4 – Programming Protocol-Independent Packet Processors, ACM CCR, July 2014

NETWORK FUNCTION VIRTUALIZATION (NFV)

NETWORK FUNCTION VIRTUALIZATION — I



- › Today's network services are highly dependent on HW appliances

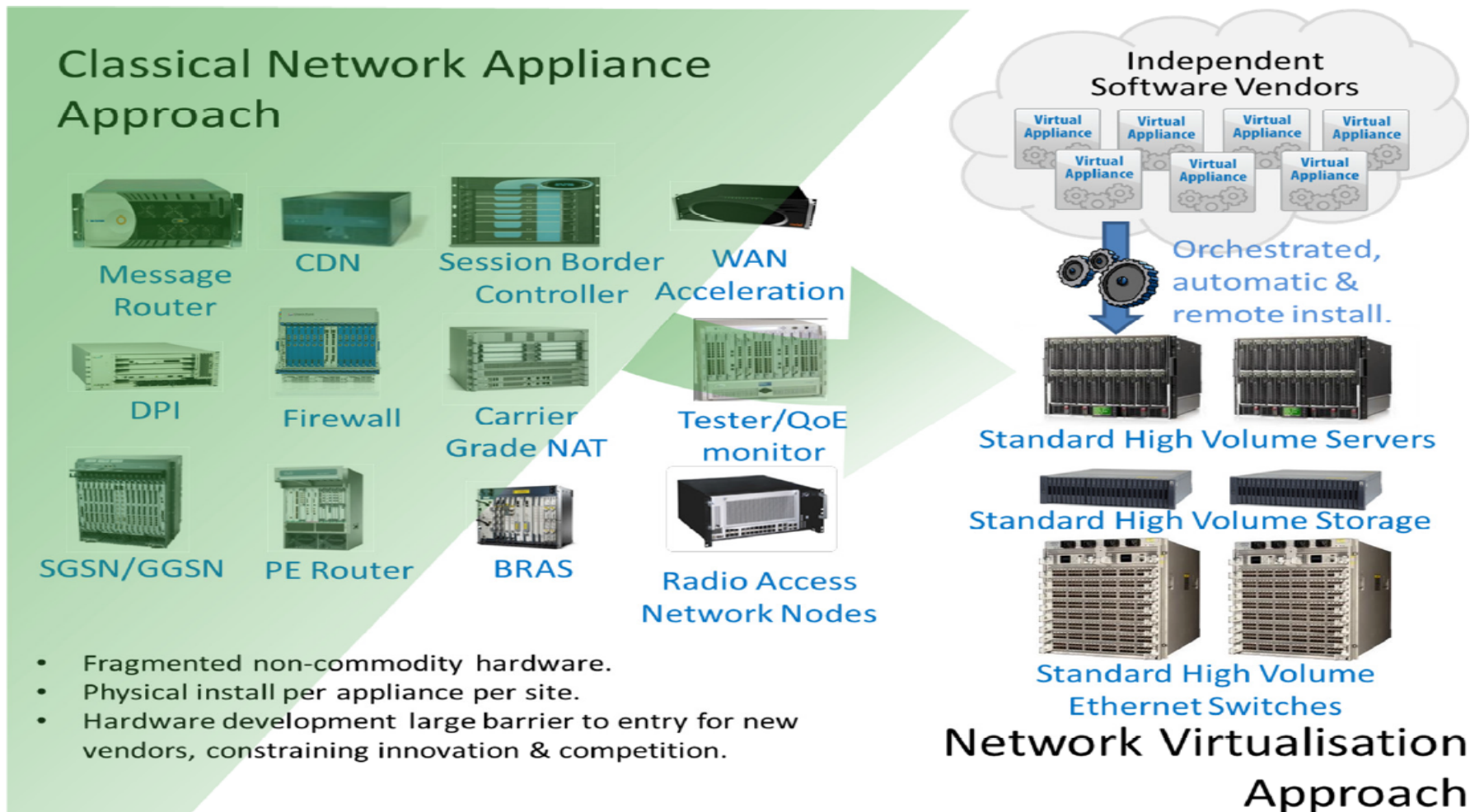


- › An end-to-end network service usually comprises several network functions (NFs) distributed over fixed locations
- › NF Examples: Firewall, Fixed and Mobile Gateways, AAA, DPI

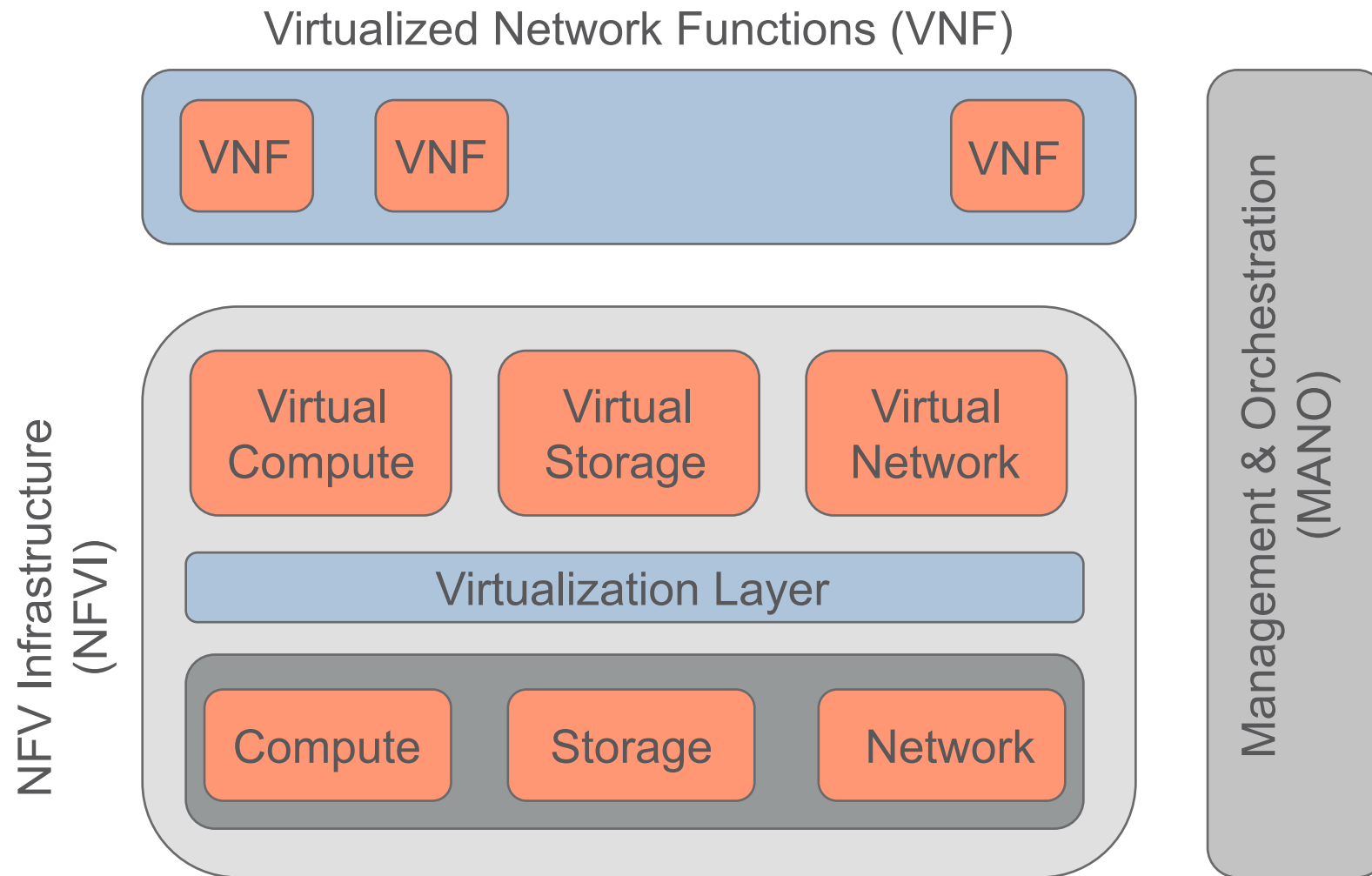
NETWORK FUNCTION VIRTUALIZATION — II



- › Virtualizing functions in SW & running them on commodity servers (e.g., in a cloud)



NFV ARCHITECTURE (ETSI VIEW)



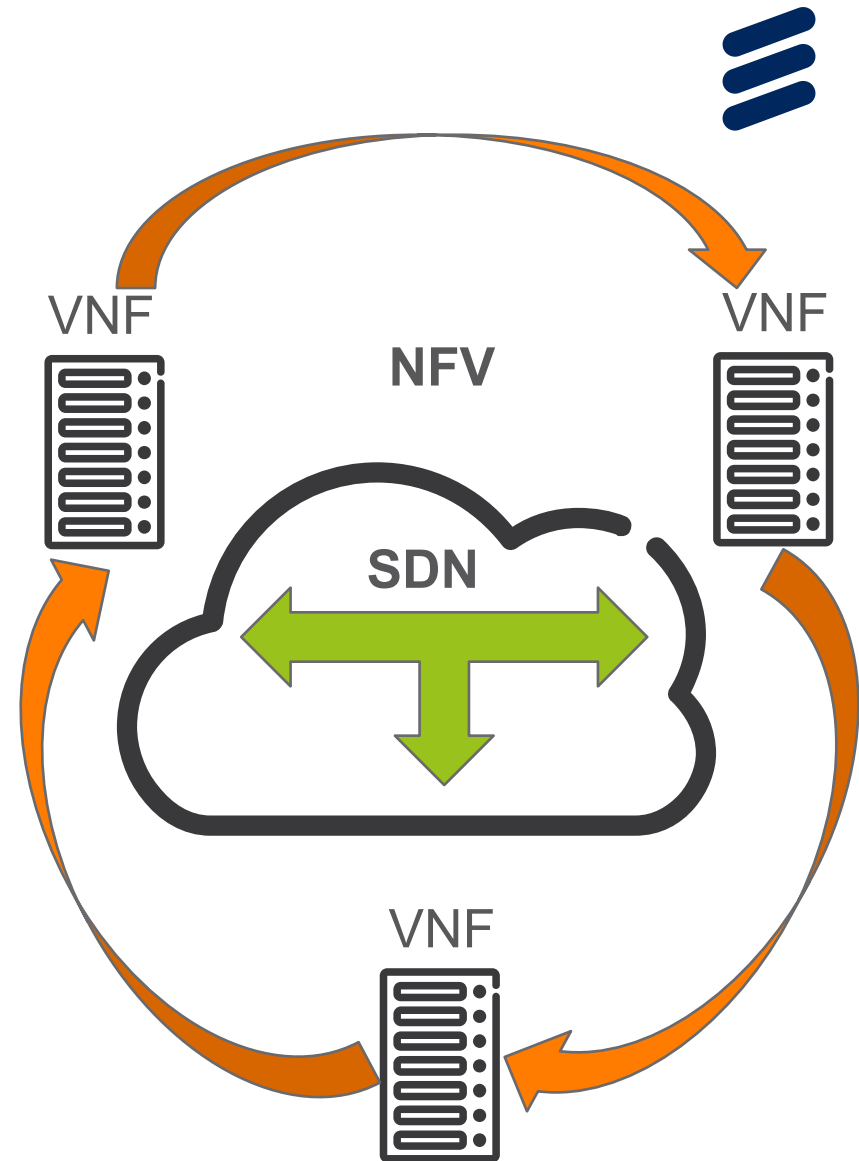
NFV BENEFITS



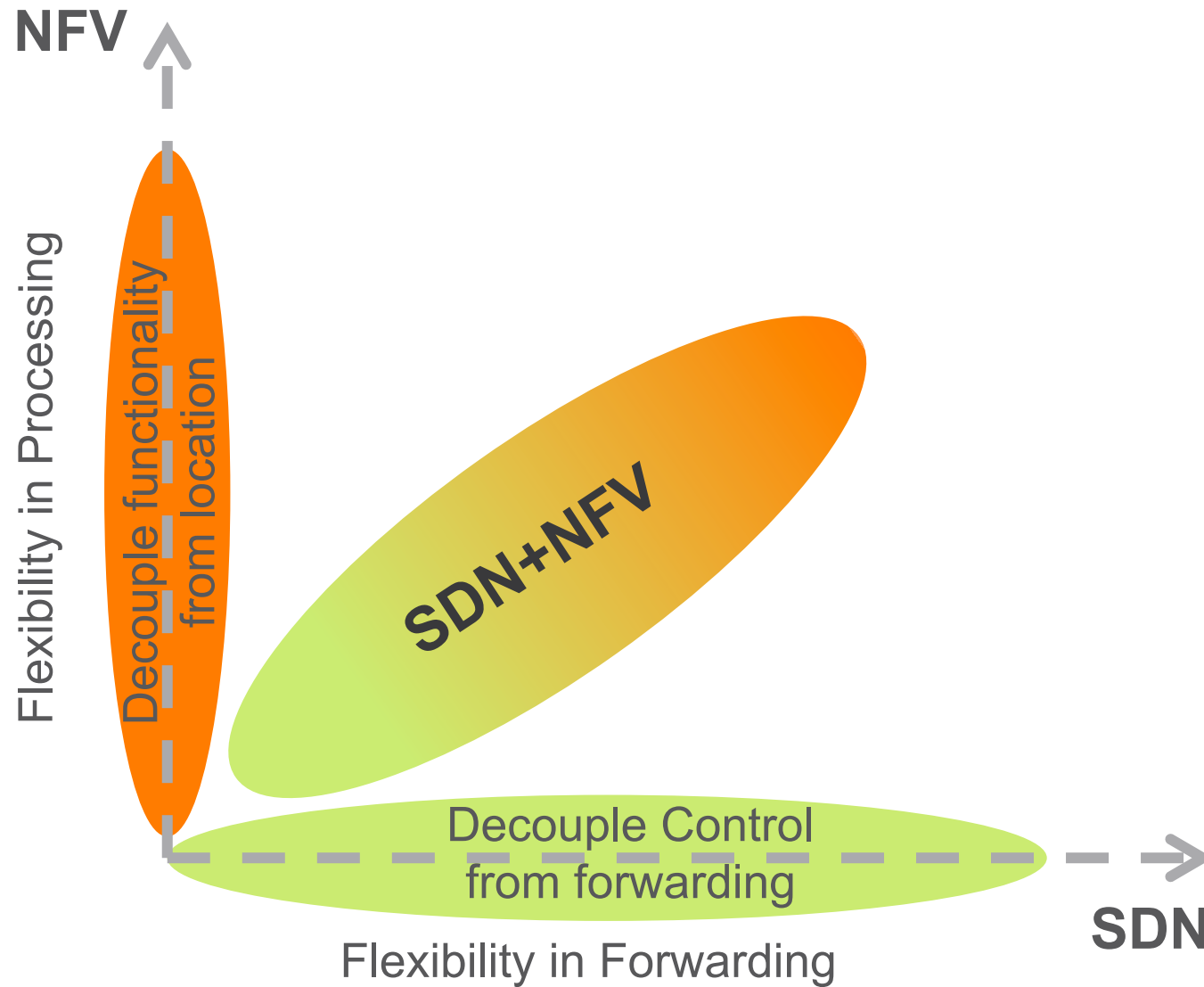
- › Efficient resource utilization
- › Agile creation of new networking services
- › Decoupling the evolution of network functions (services) from underlying platforms
- › Reduced Cost (CapEX & OpEX)

SDN VS. NFV

- › SDN → flexible forwarding & steering of traffic in a physical or virtual network environment
- › NFV → flexible placement of virtualized network functions across the network & cloud
- › SDN & NFV are complementary tools for achieving full network programmability



FLEXIBILITY WITH SDN & NFV



SDN & NFV IN
ACTION:
UNIFYING THE
PROGRAMMABILITY OF
CLOUD AND CARRIER
INFRASTRUCTURE

UNIFY CONSORTIUM

2013 NOVEMBER - 2016 APRIL



Major Service Providers:



Research Institutes:



Major Vendors:



Universities:



Project Management: **eict**

UNIFY REQUIREMENTS



› *Today, rigid network control limits the flexibility of service creation*

› Operators



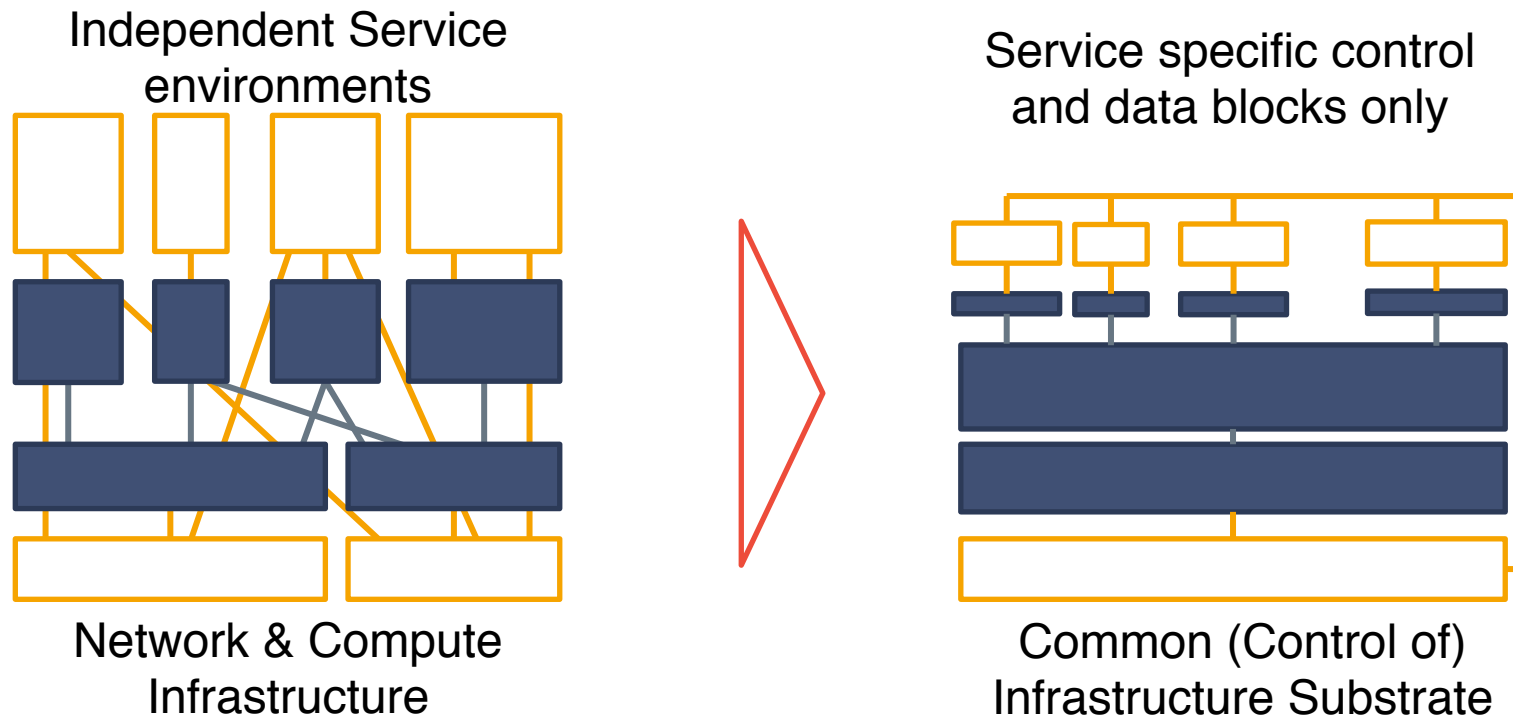
- › Agility/Velocity
- › Flexibility/Granularity
- › Simplicity/Automation
- › Scalability
- › Lower OPEX
- › Integration/Shareconomy

› Users



- › Rich services,
- › Quality of experience,
- › Rapid provisioning of elastic services,
- › On-demand SLA configuration & monitoring,
- › “Follow-me” services

UNIFIED PRODUCTION ENVIRONMENT



Unified Production Environment with dynamic service creation platform, leveraging a fine-granular service chaining architecture.

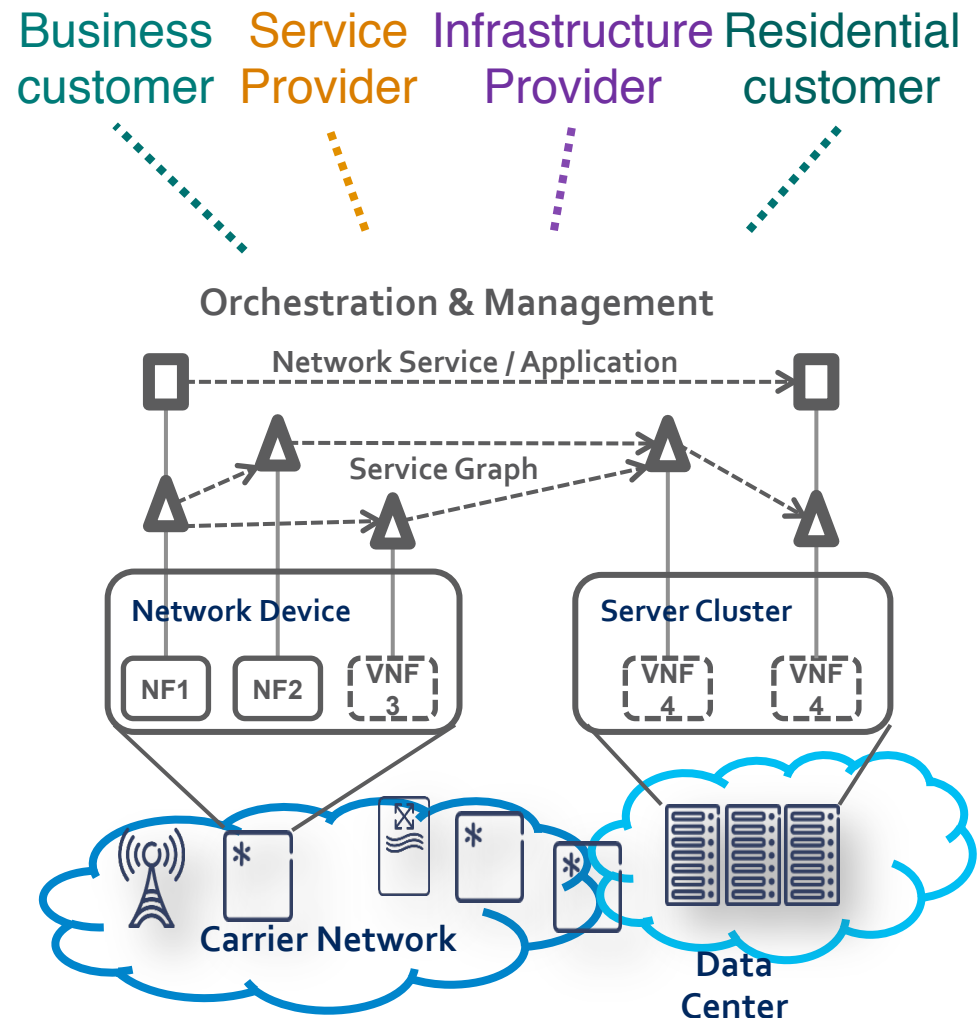
UNIFY FOCUS SEAMLESS INTEGRATION OF CLOUD AND NETWORK



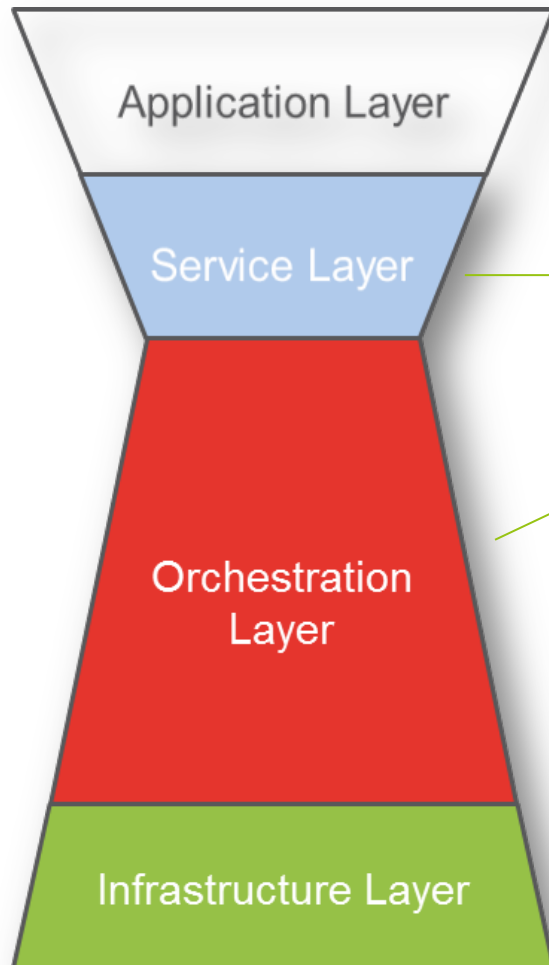
Invocation of
Dynamic Service Chains
(Programmability)
UNIFY Control Plane

Joint orchestration in Network &
Cloud (Abstractions)
UNIFY Control Plane

Data performance
optimized infrastructure
virtualization
(x86 based architecture)
UNIFY Universal Node



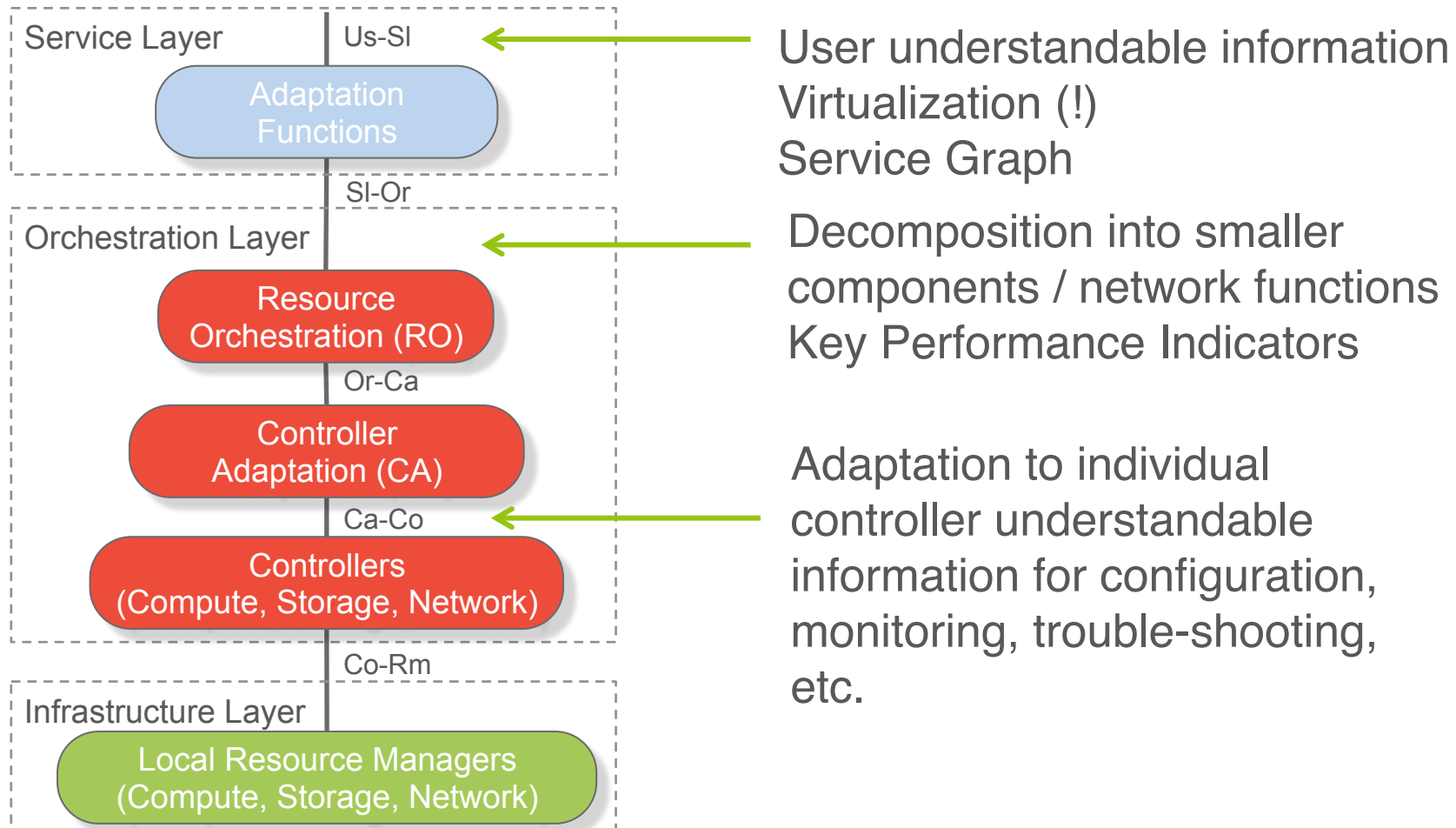
UNIFY LAYERED ARCHITECTURE



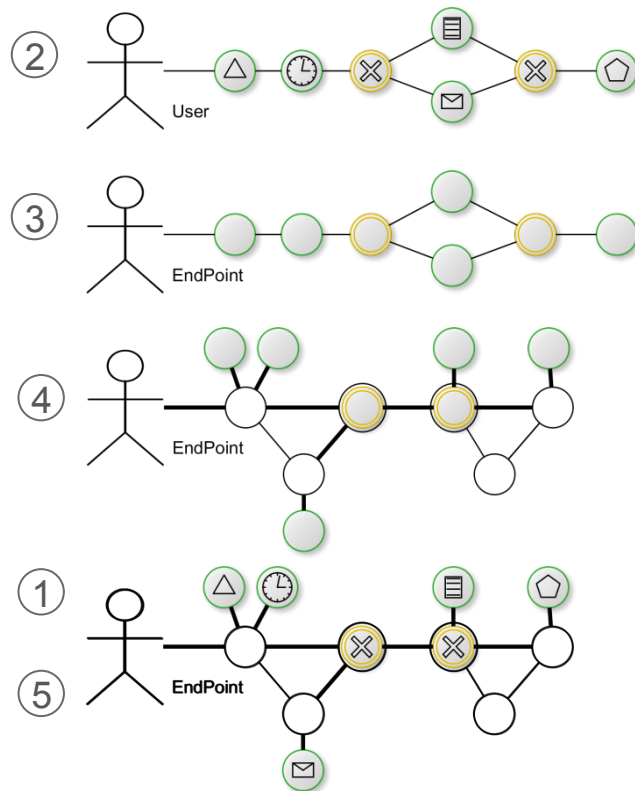
UNIFY Control Plane

- **Services** decomposition and abstraction
- **Orchestration** for Network Function Forwarding Graphs (NF-FG)
- **Combined Compute, Storage & Network** abstraction over all resources
 - forwarding elements,
 - compute host capabilities,
 - hardware based network function capabilities,
 - data plane links

UNIFY SLICING THE ELEPHANT – SEPARATION OF PROBLEMS

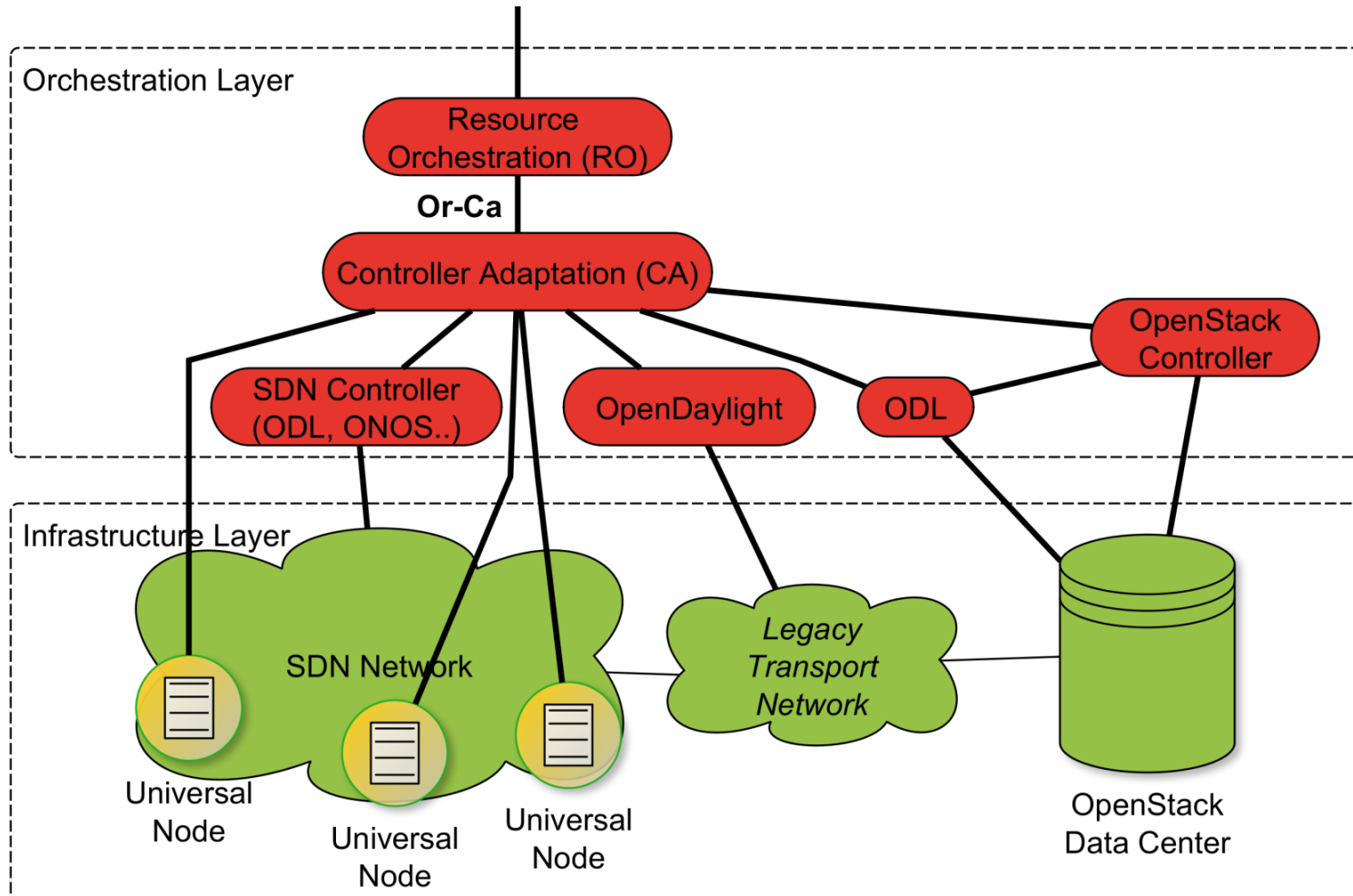


UNIFY ABSTRACTION – AN EXAMPLE

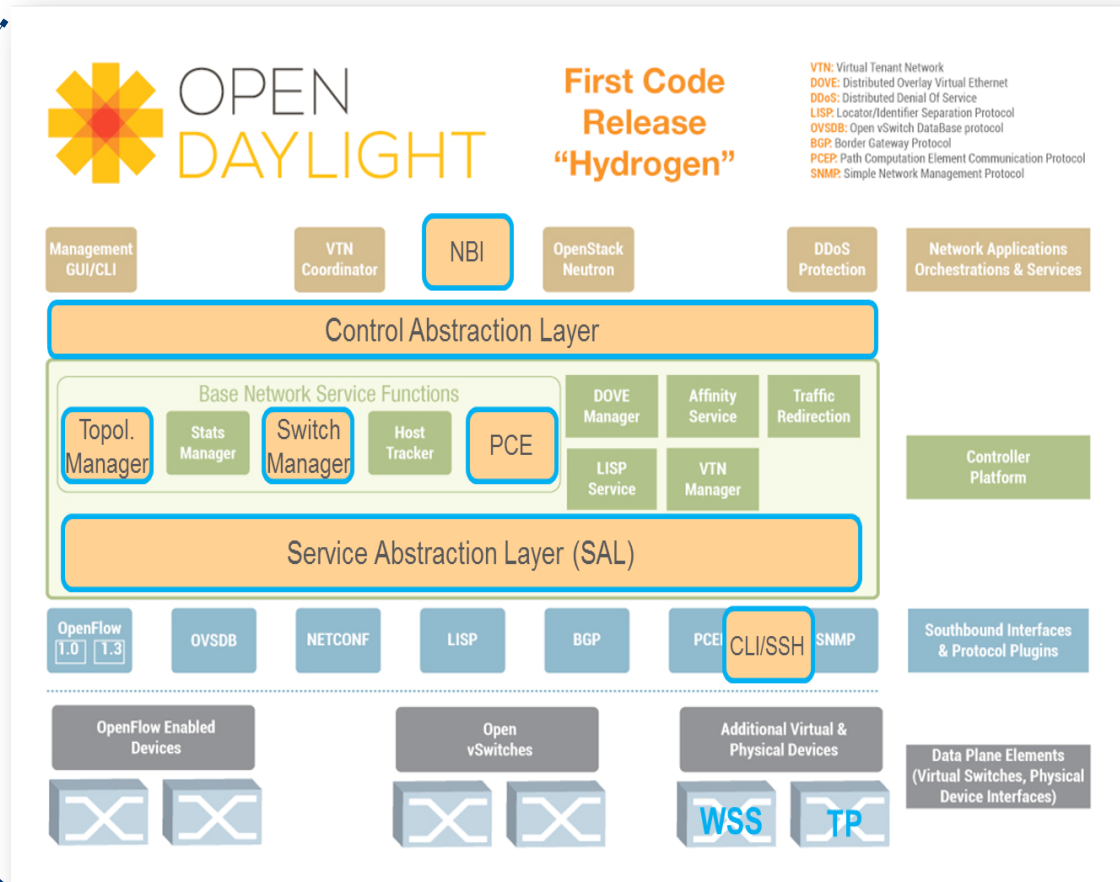
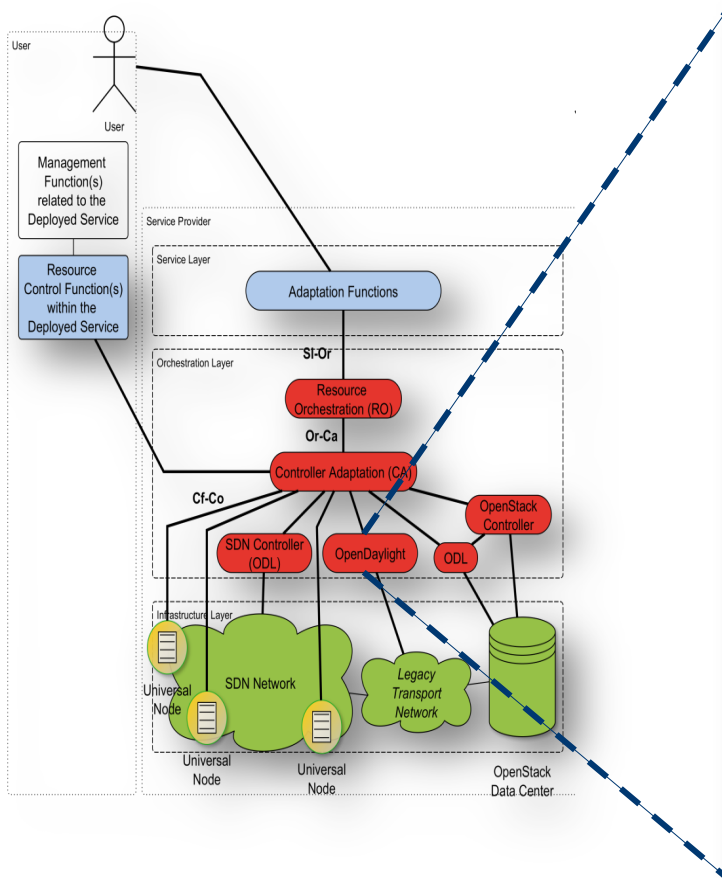


- Service Graph
- Abstract *Network Function Forwarding Graph*
- Abstract Resource Mapping
- Physical Infrastructure
 - Compute, Storage, Network and Topology
 - **Instantiated Service Graph**

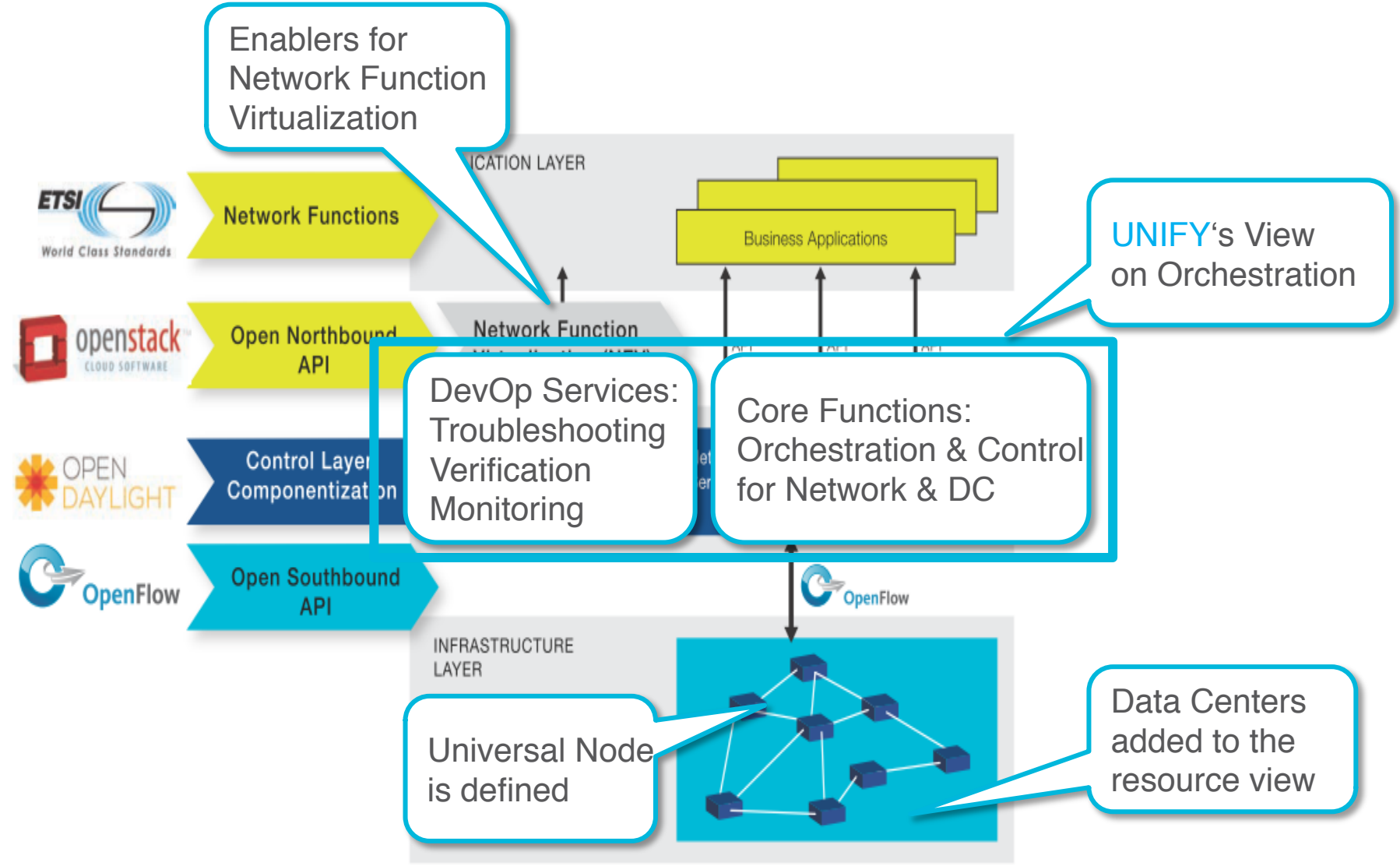
UNIFY EXAMPLE FOR REALIZATION



UNIFY IMPLEMENTATION OPTION – LEGACY OPTICAL TRANSPORT



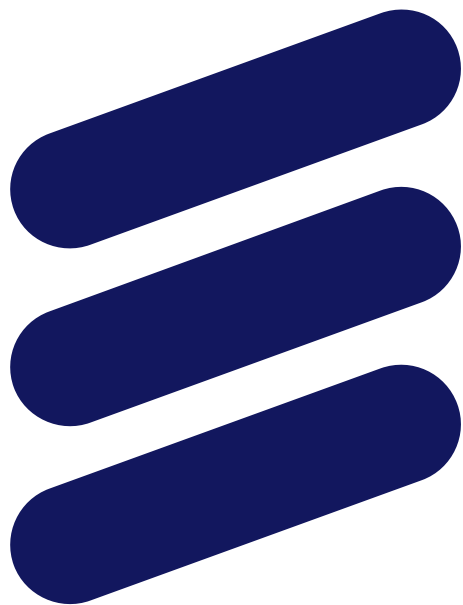
UNIFY VS. NFV AND SDN INDUSTRY



LINKS TO RELEVANT BODIES



- › Open Networking Foundation (ONF)
 - › www.opennetworking.org
- › ETSI NFV Initiative
 - › <http://www.etsi.org/technologies-clusters/technologies/nfv>
- › OpenDayLight Project
 - › <http://www.opendaylight.org>
- › IETF ForCES WG
 - › <http://tools.ietf.org/wg/forces>
- › IETF Service Function Chaining WG
 - › <https://datatracker.ietf.org/wg/sfc/charter>
- › IRTF SDN Research Group
 - › <https://irtf.org/sdnrg>
- › SDN in IEEE
 - › <http://sdn.ieee.org>



ERICSSON